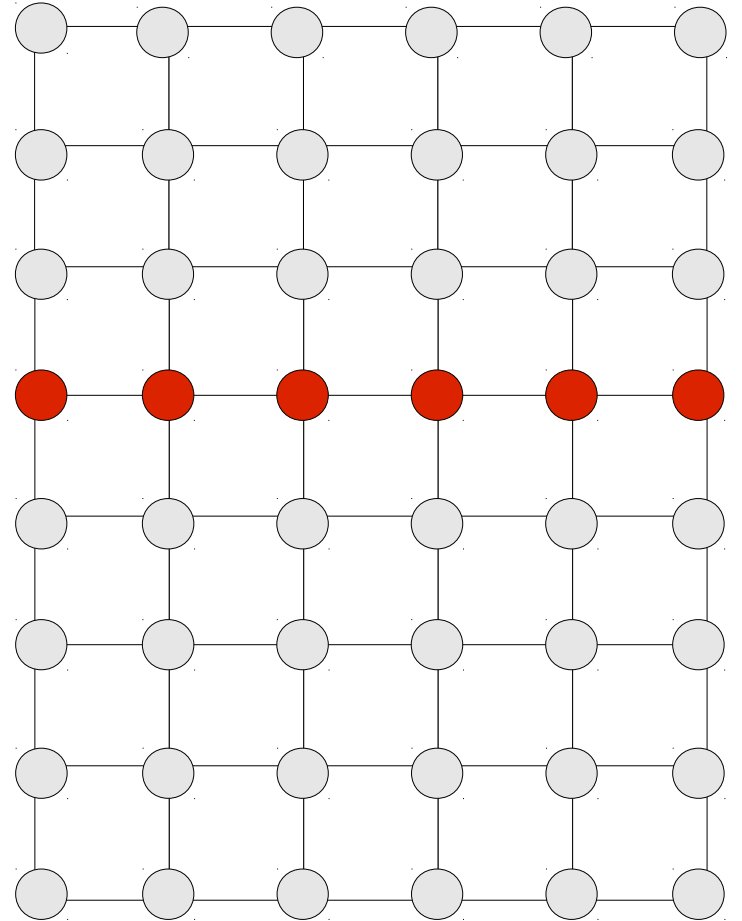

The Planar Separator Theorem

by Jan Dreier

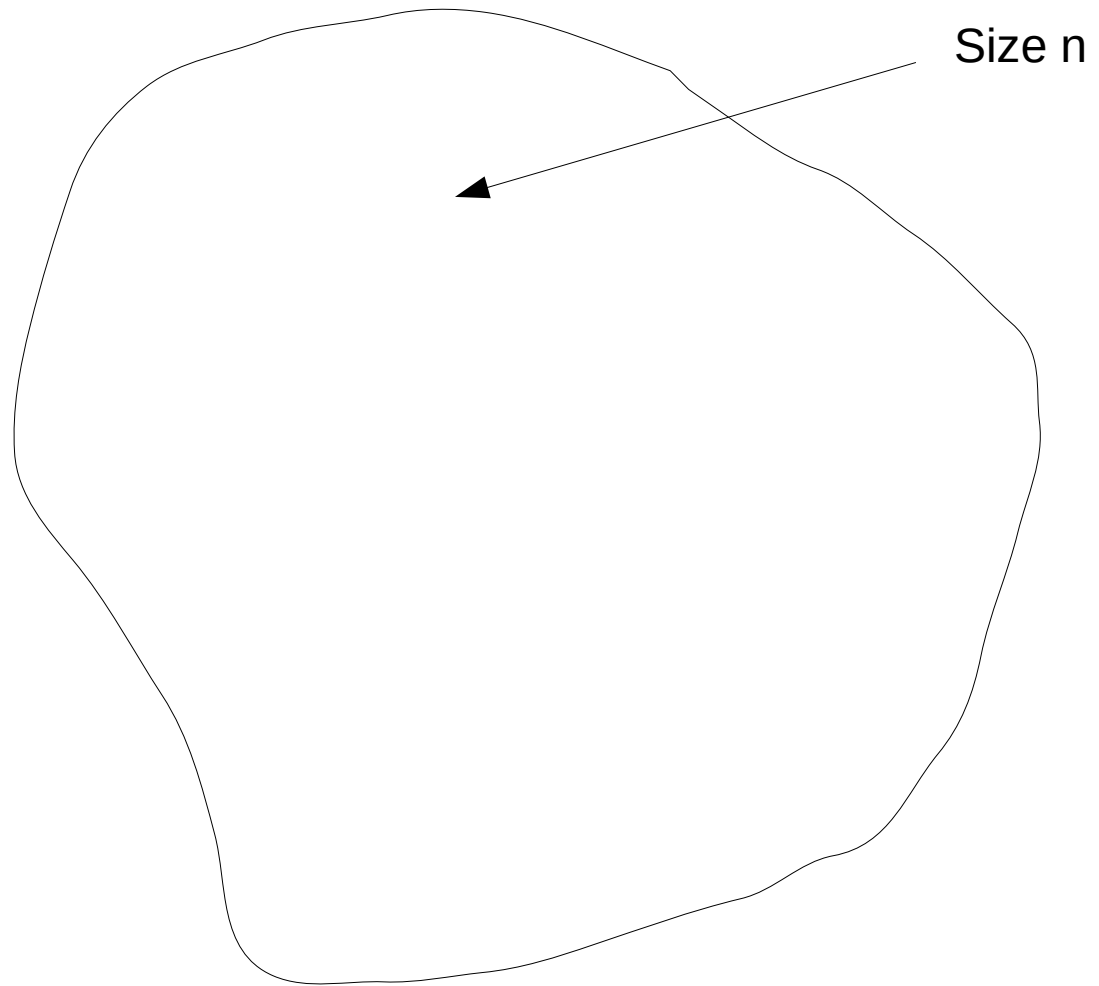
The Planar Separator Theorem

- **Separator:** A set of vertices, that when removed divide the graph in two non-connected components.
- **Goals:**
 - Equally-sized components
 - Small Separator
- **Applications:**
 - “Divide-and-conquer” Algorithms



- **Application-Example:** Approximation of Maximum Independent Set

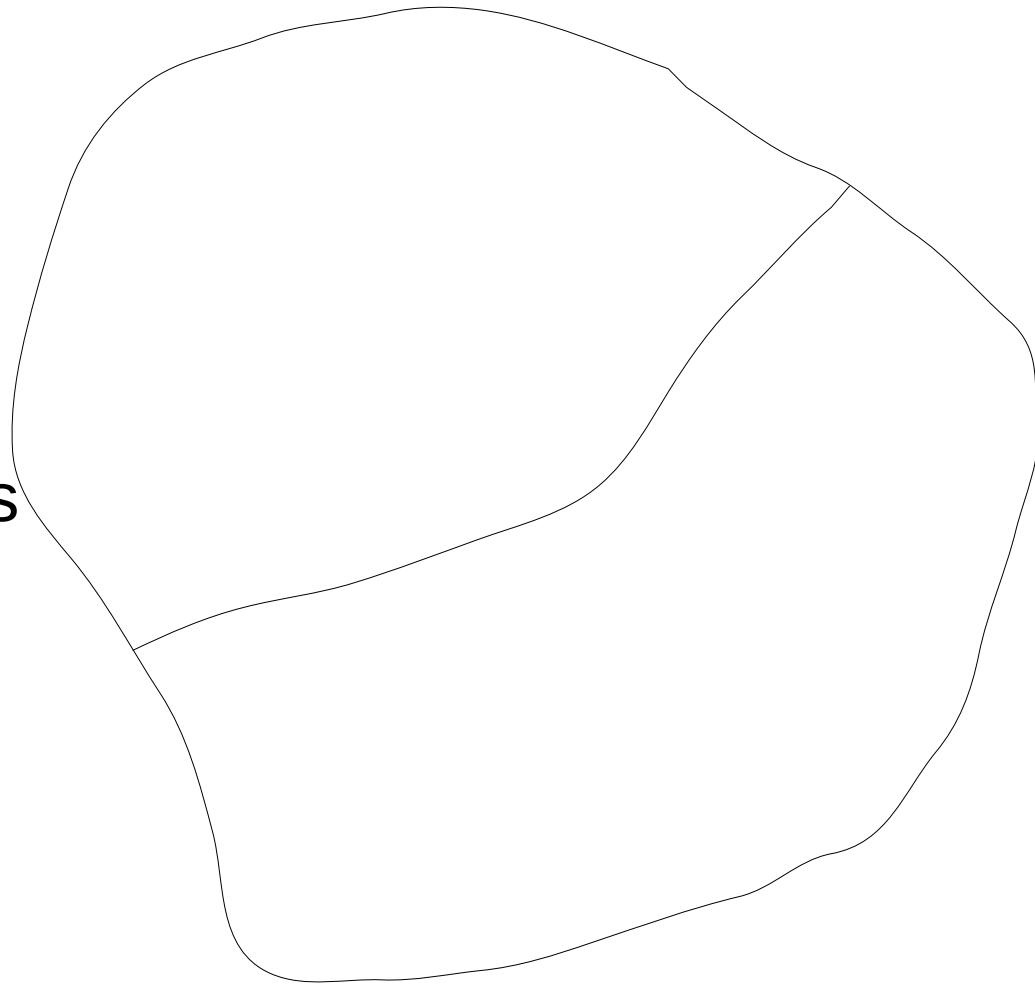
Problem is NP-complete



- **Application-Example:** Approximation of Maximum Independent Set

Problem is NP-complete

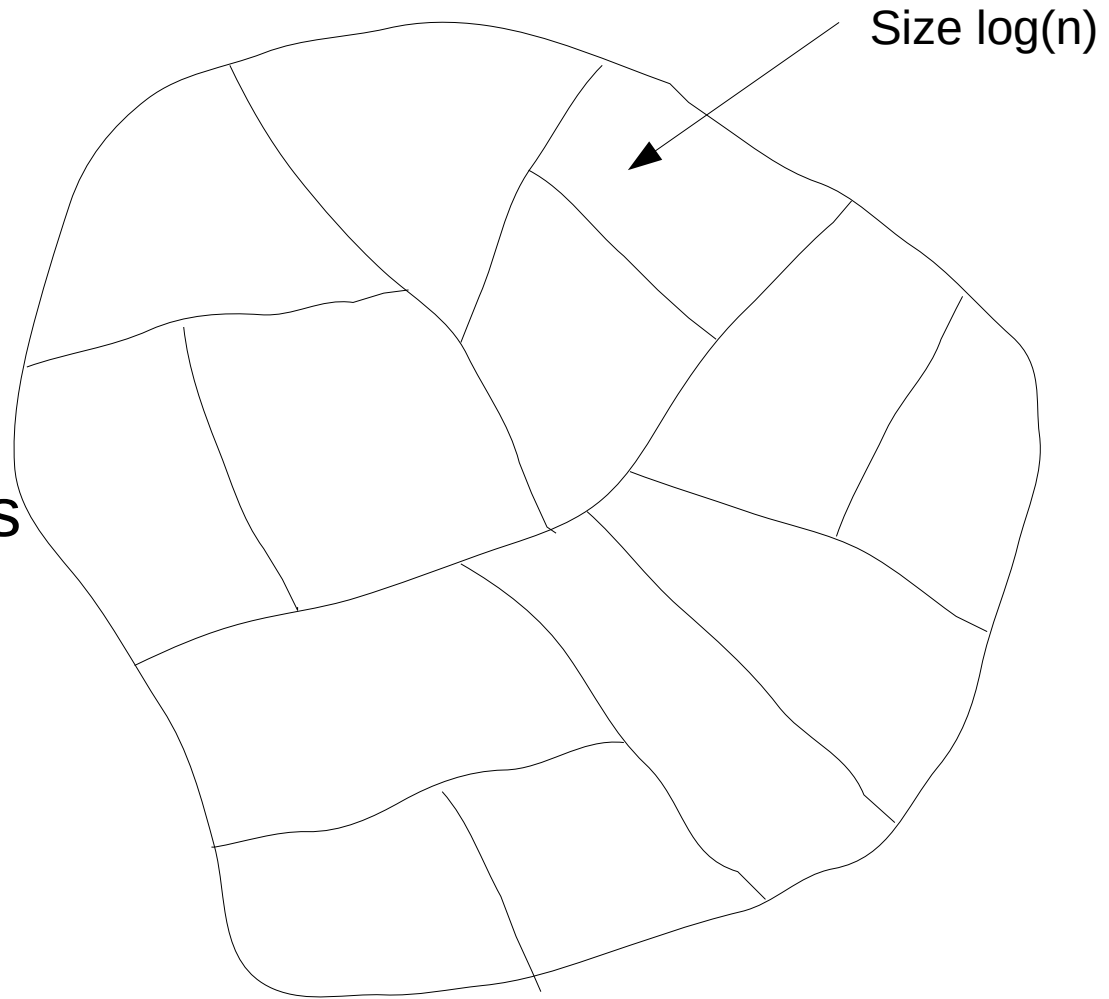
Separate into small connected components



- **Application-Example:** Approximation of Maximum Independent Set

Problem is NP-complete

Separate into small connected components



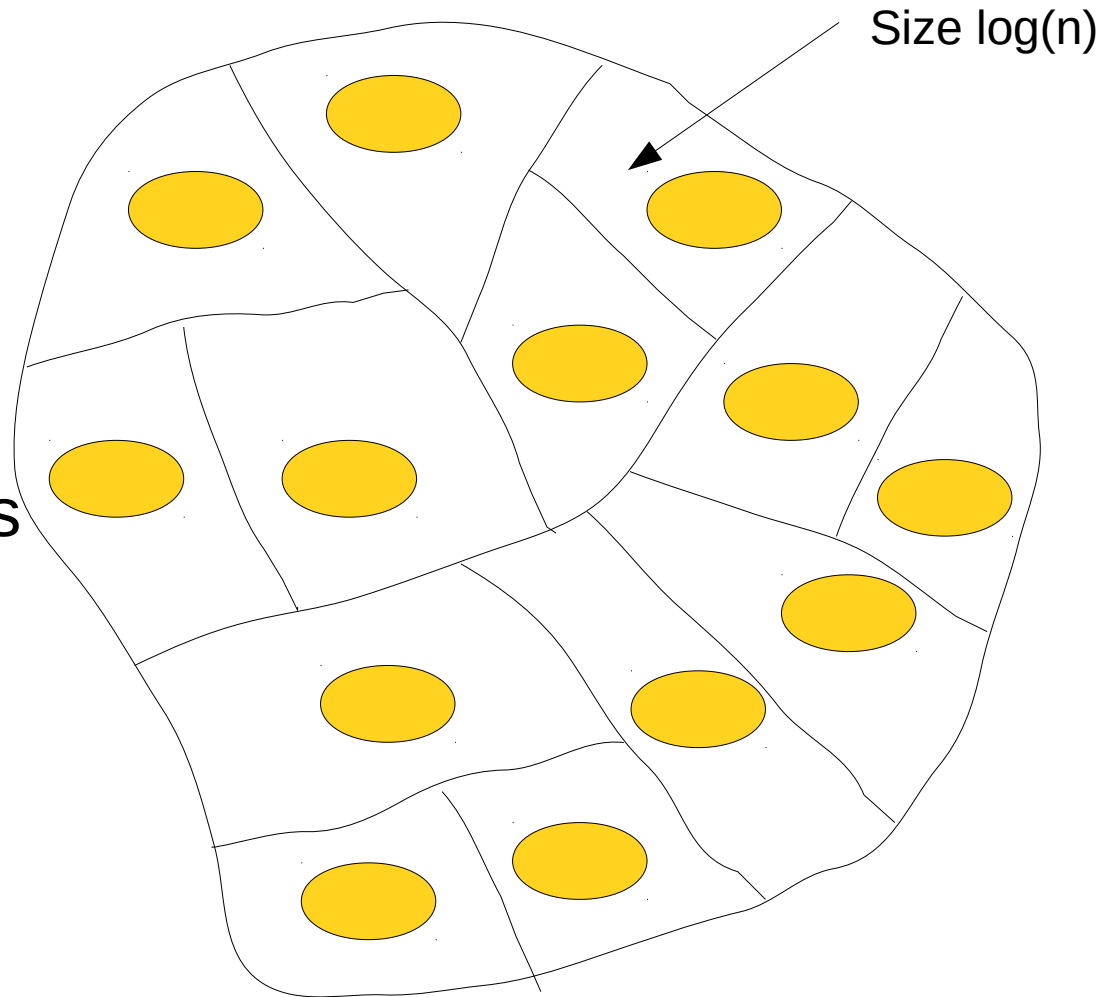
The Planar Separator Theorem

- **Application-Example:** Approximation of Maximum Independent Set

Problem is NP-complete

Separate into small connected components

Find Maximum Independent Set in each component



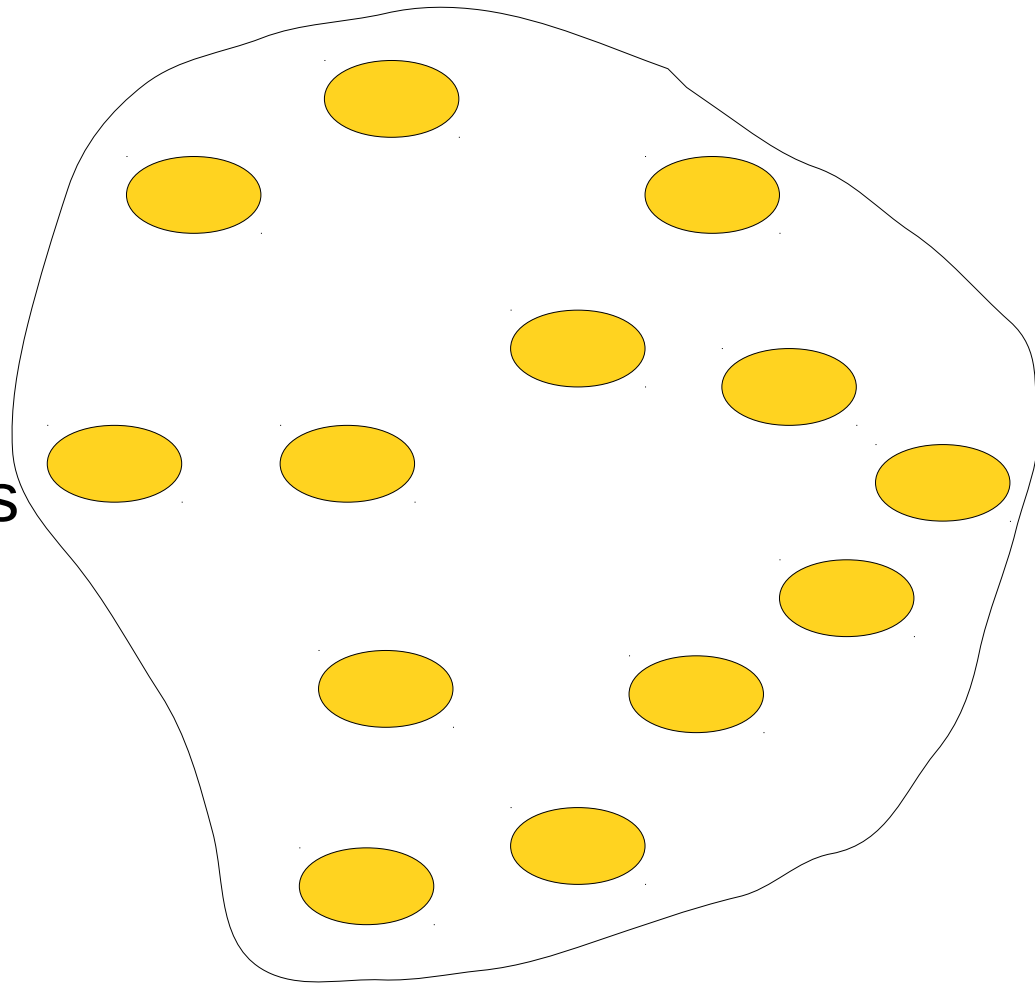
- **Application-Example:** Approximation of Maximum Independent Set

Problem is NP-complete

Separate into small connected components

Find Maximum Independent Set in each component

Return union of all Independent Sets

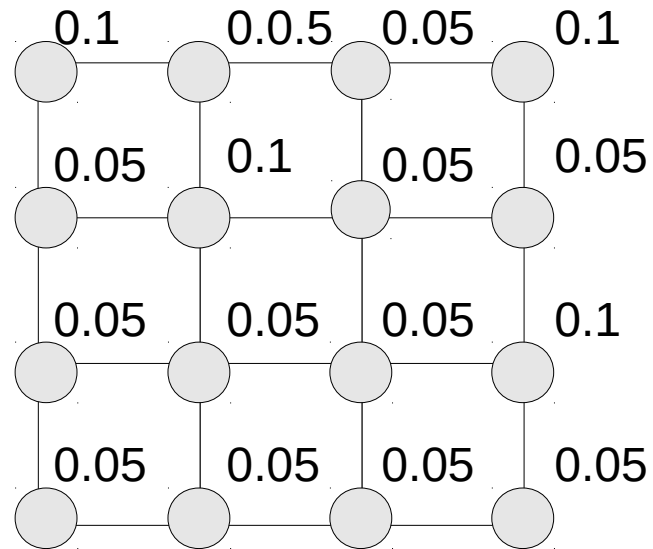


- **Planar Separator Theorem (Lipton and Tarjan, 1979):**
 - For a weighted planar graph with n nodes and a summed weight of no more than 1, there is a separator of size $2\sqrt{2}\sqrt{n}$ dividing the graph in two components with weight less than $\frac{2}{3}$.

- **Planar Separator Theorem (Lipton and Tarjan, 1979):**

- **Example:**

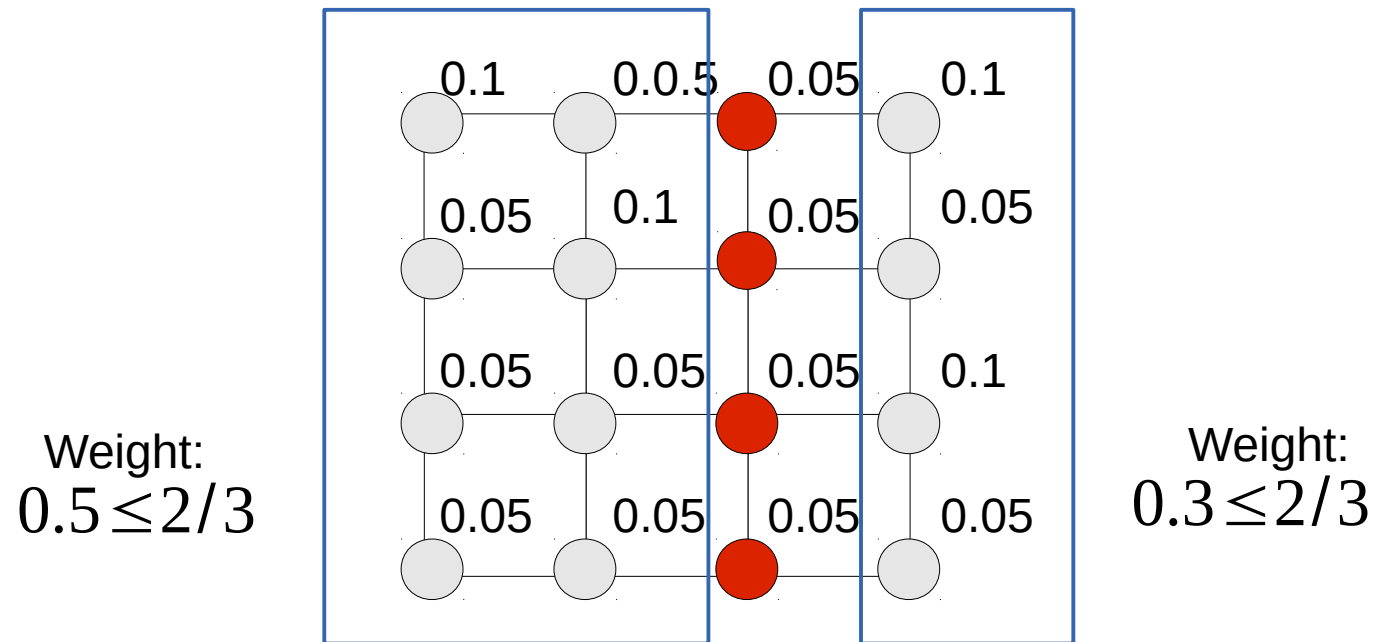
- planar graph
- n nodes
- weight of 1



- **Planar Separator Theorem (Lipton and Tarjan, 1979):**

- **Example:**

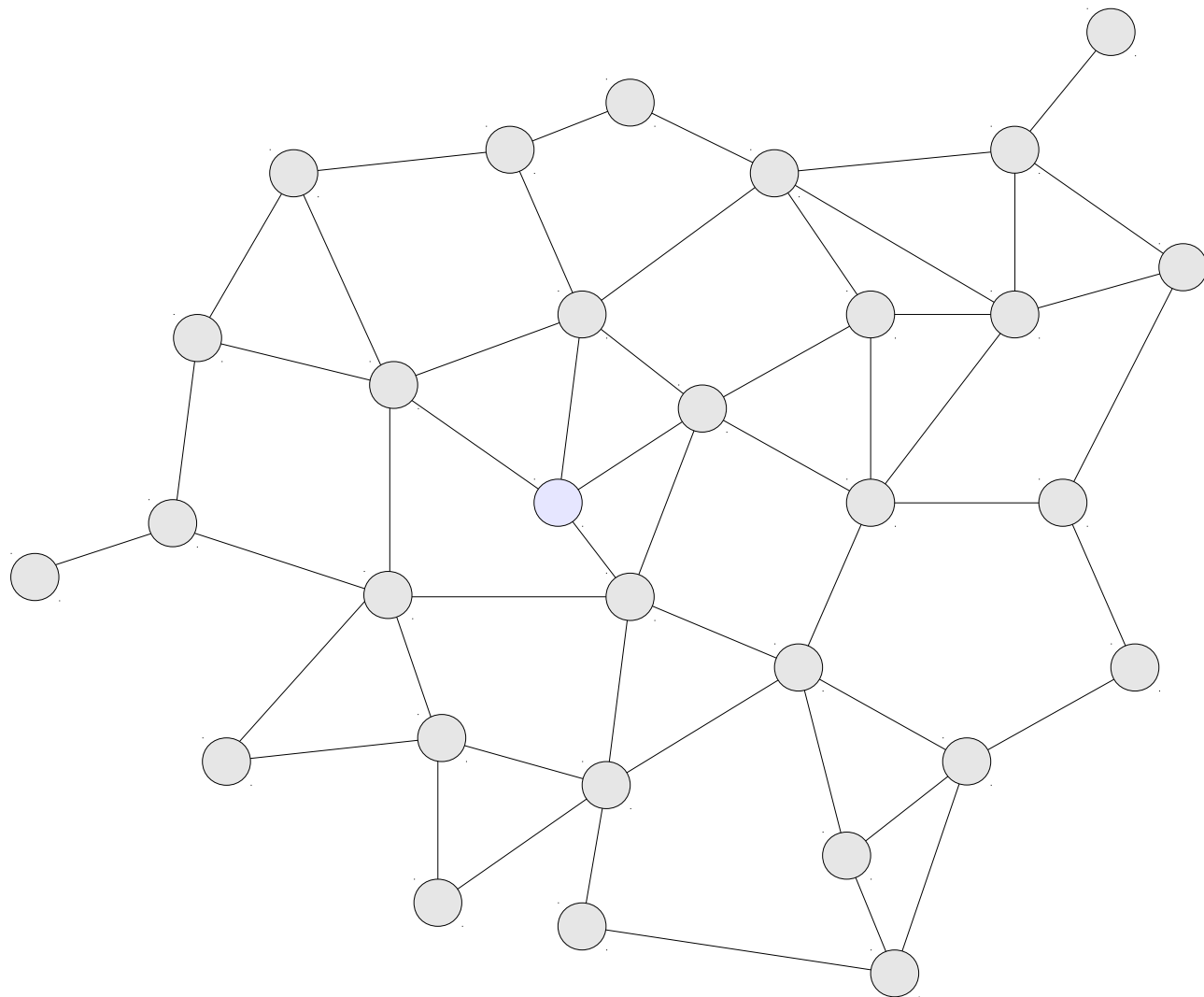
- planar graph
- n nodes
- weight of 1



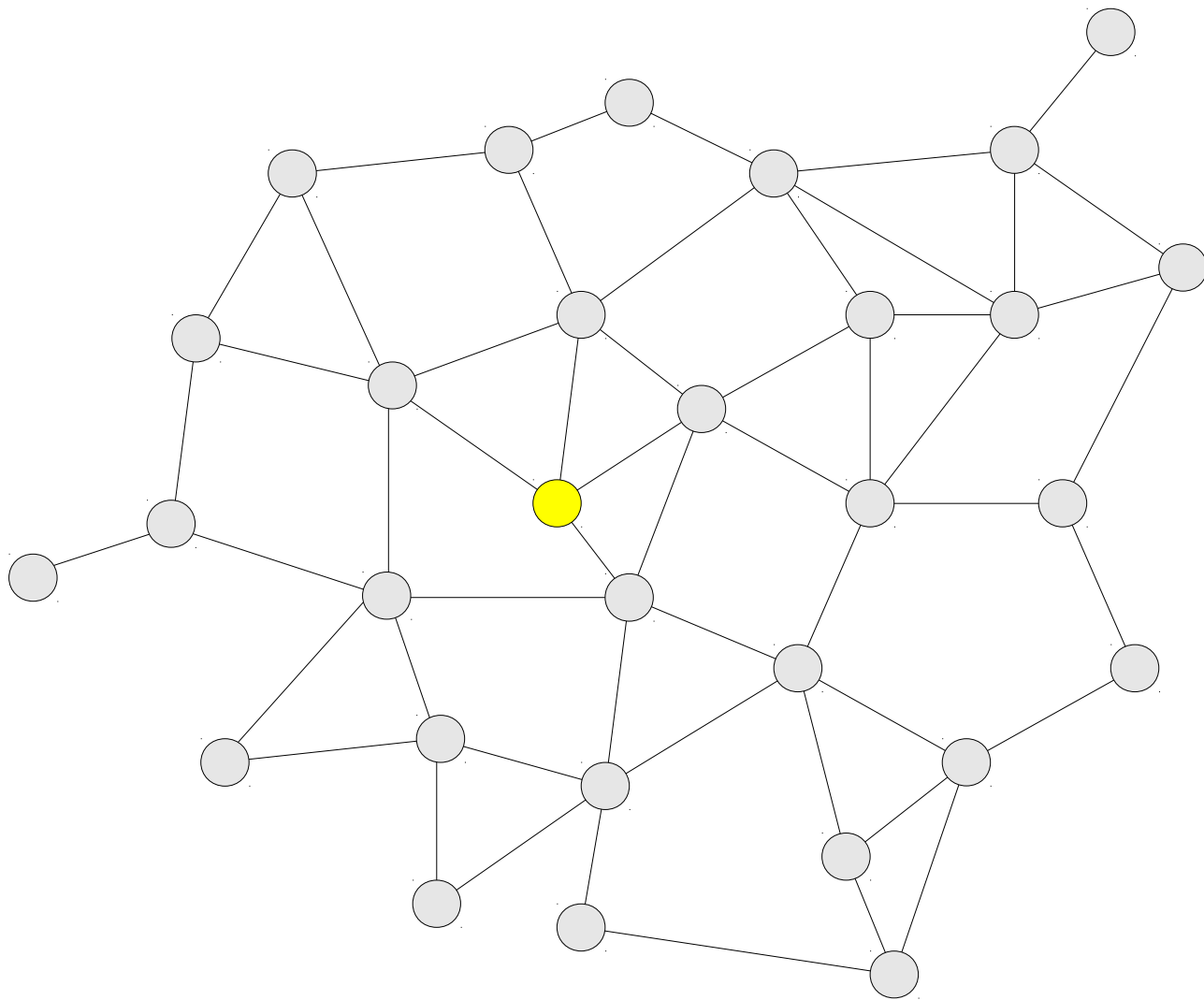
Separator: $4 \leq 2 \sqrt{(2)} \sqrt{(16)}$

- **How to find a separator in a graph?**

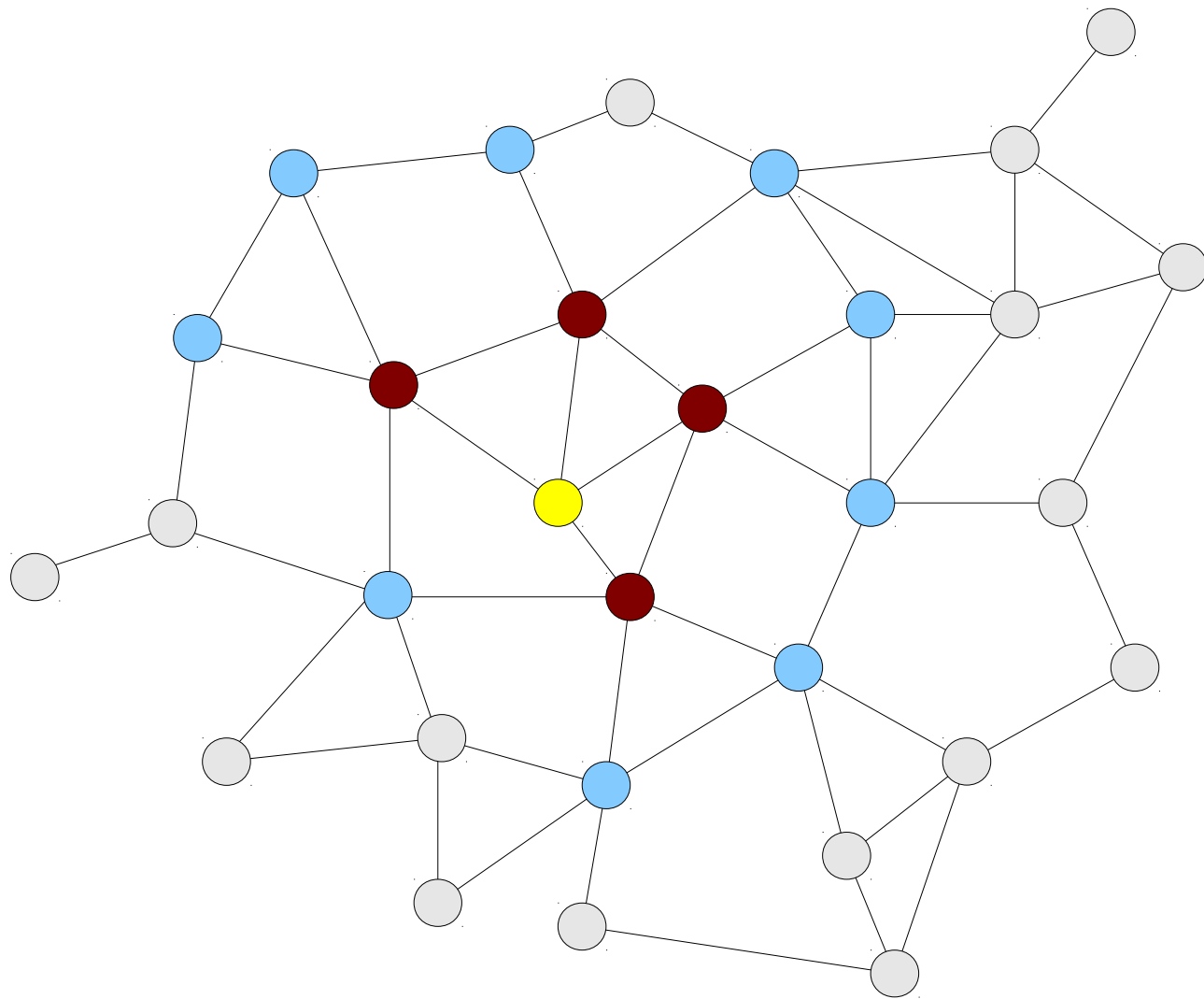
Levels



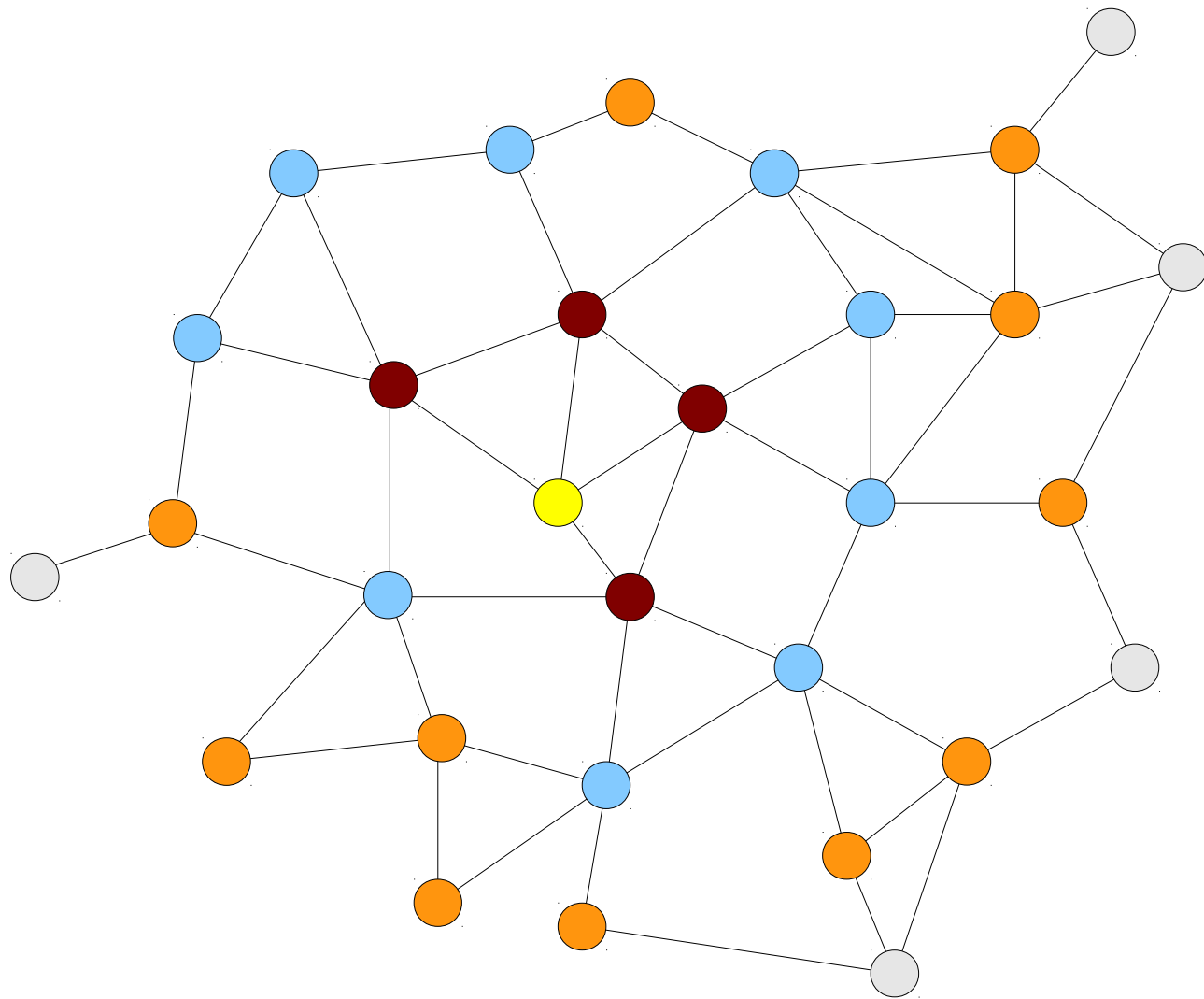
Levels



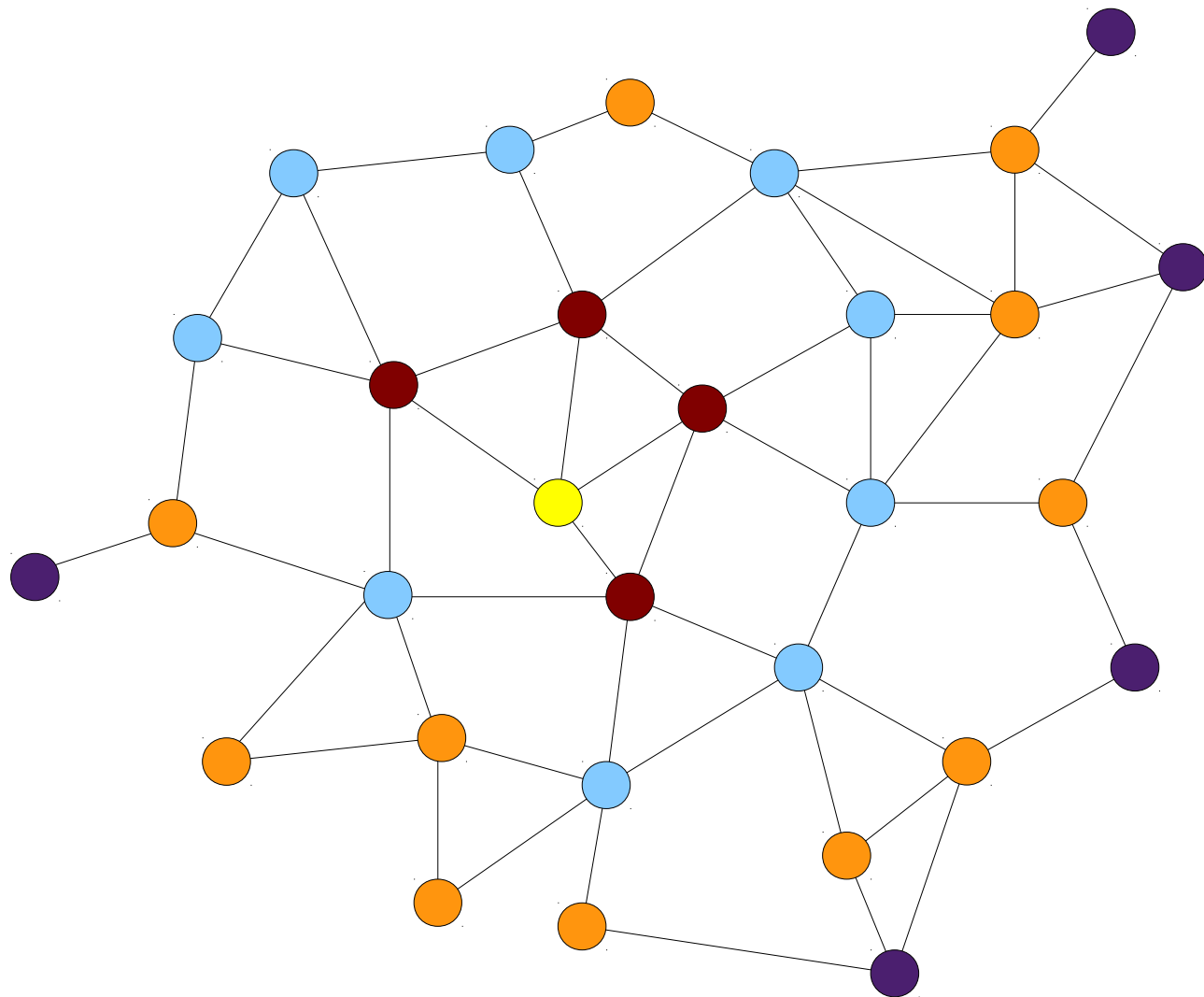
Levels



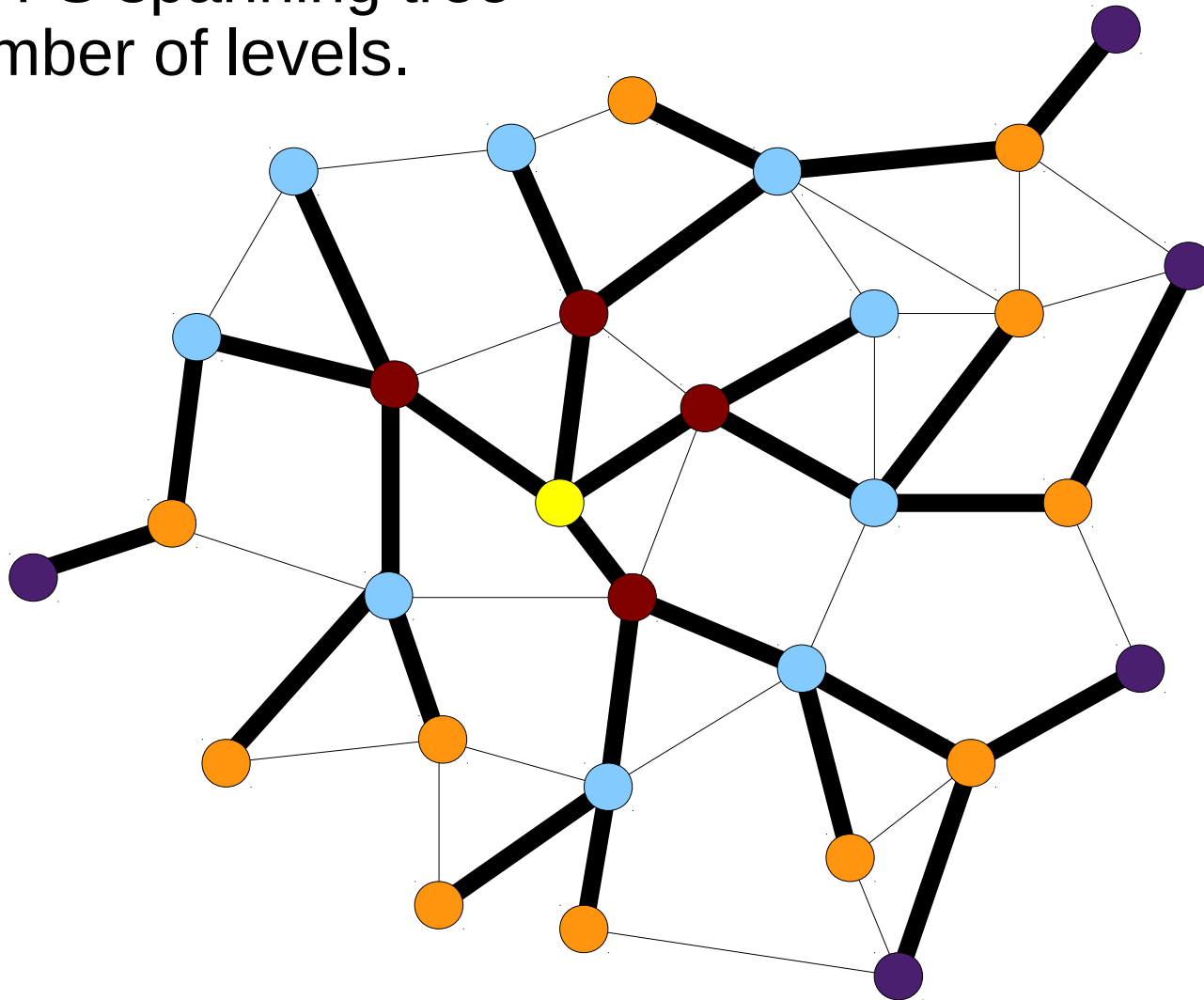
Levels



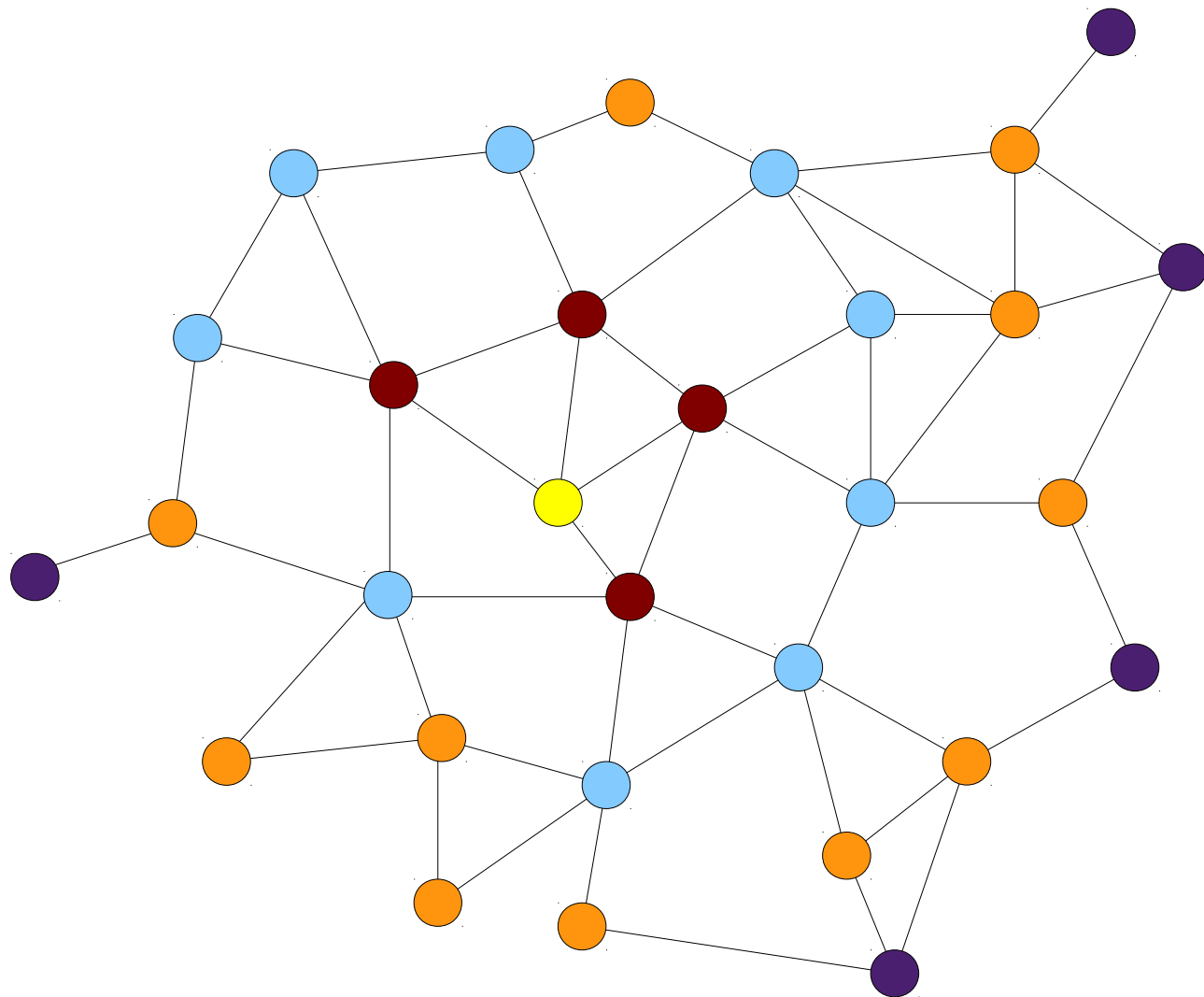
Levels



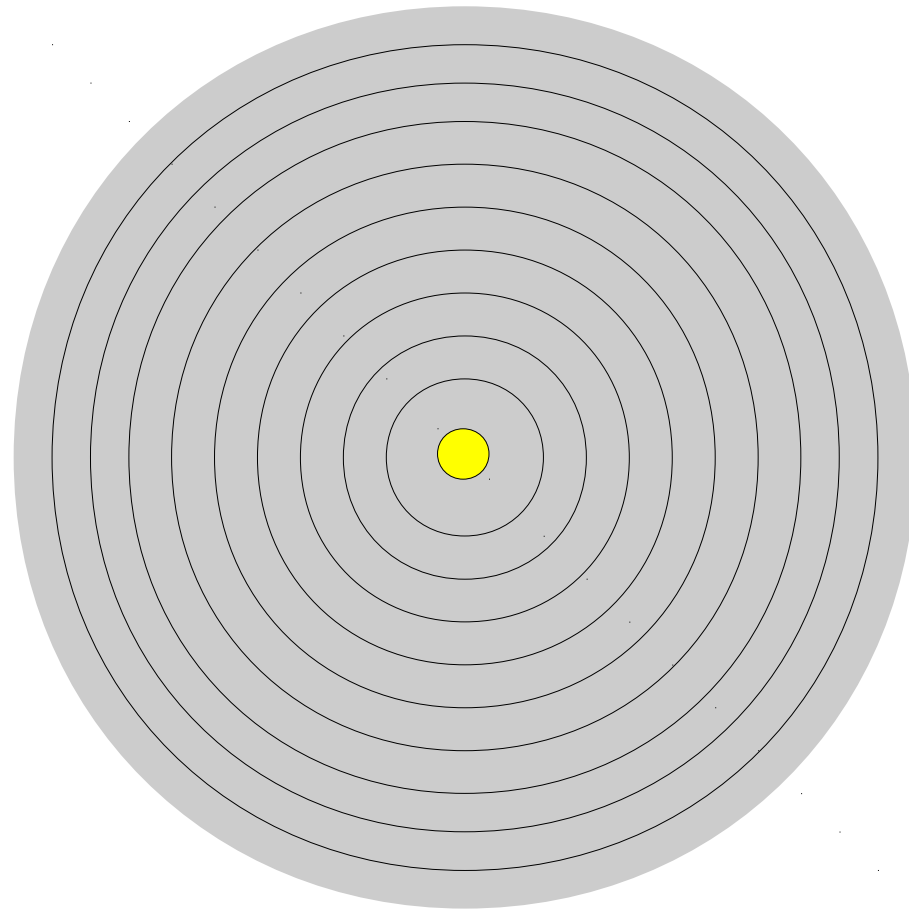
First observation:
Depth of DFS spanning tree
equals number of levels.

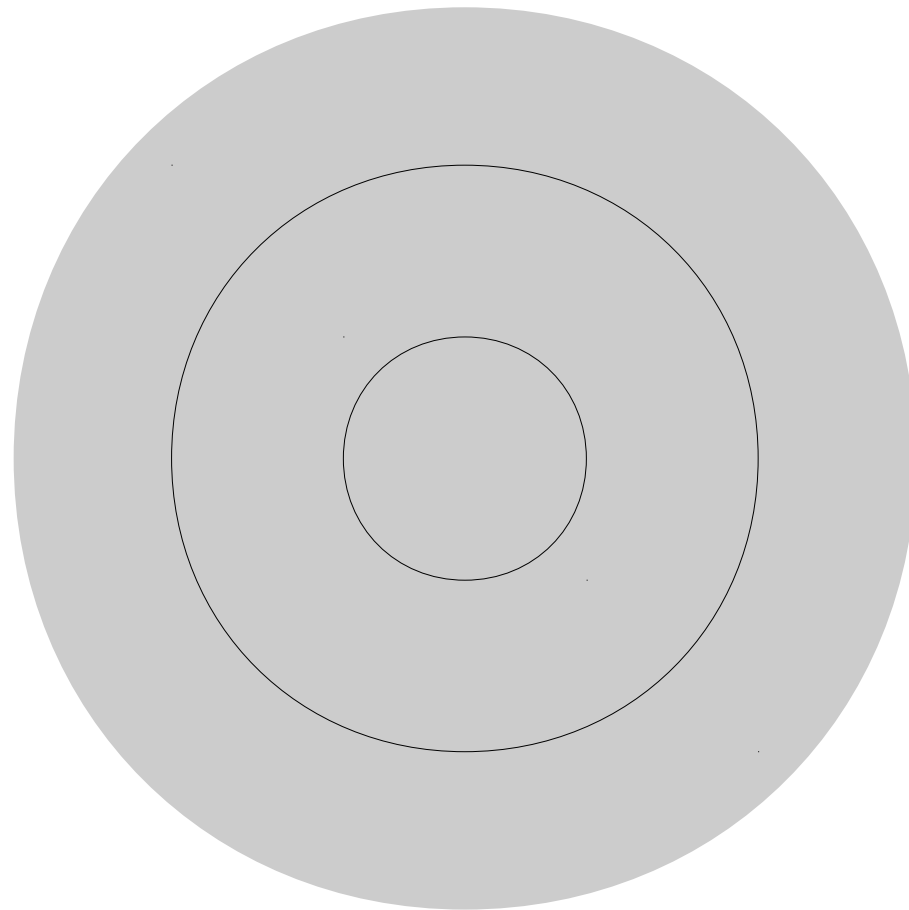


Levels

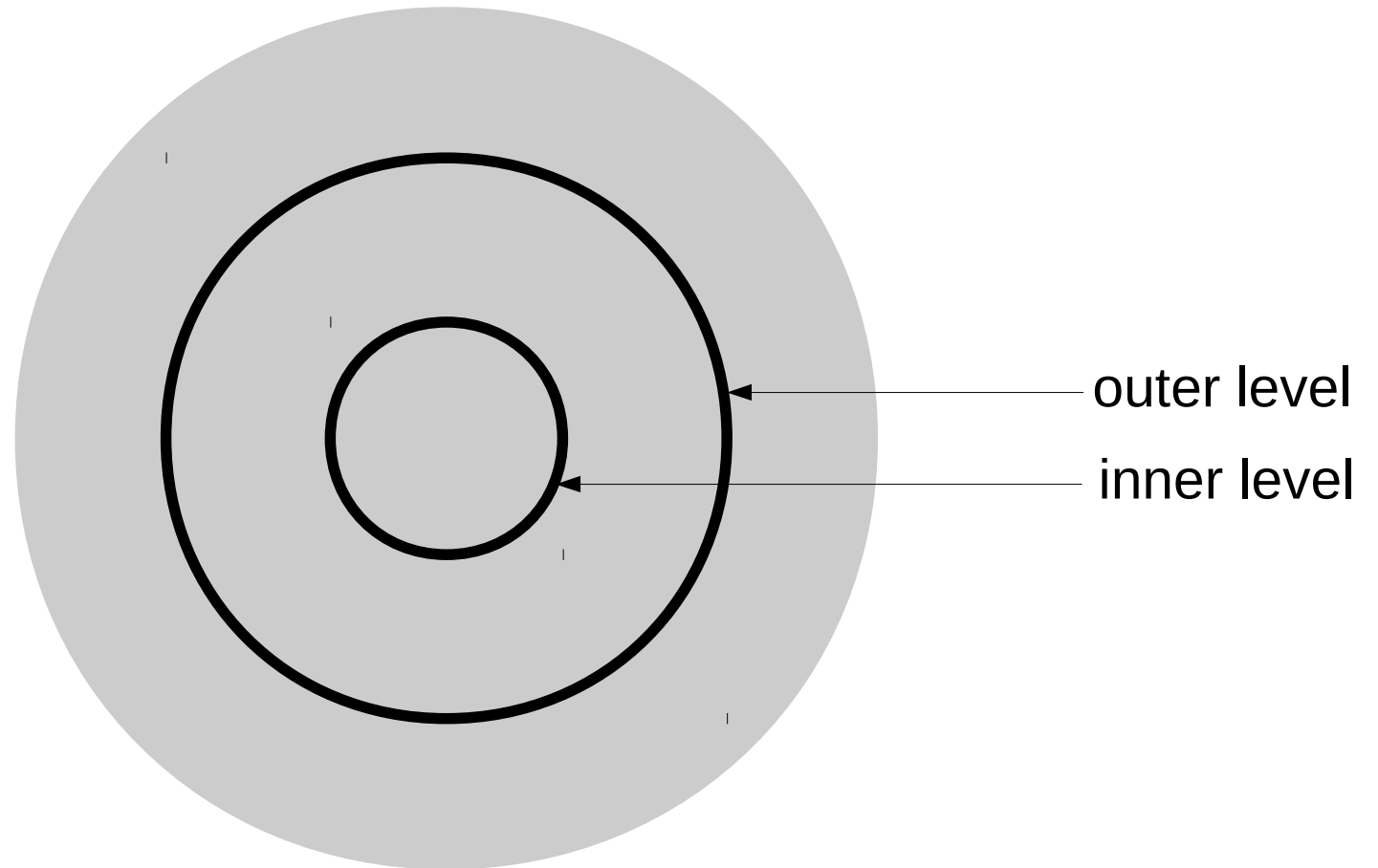


Levels

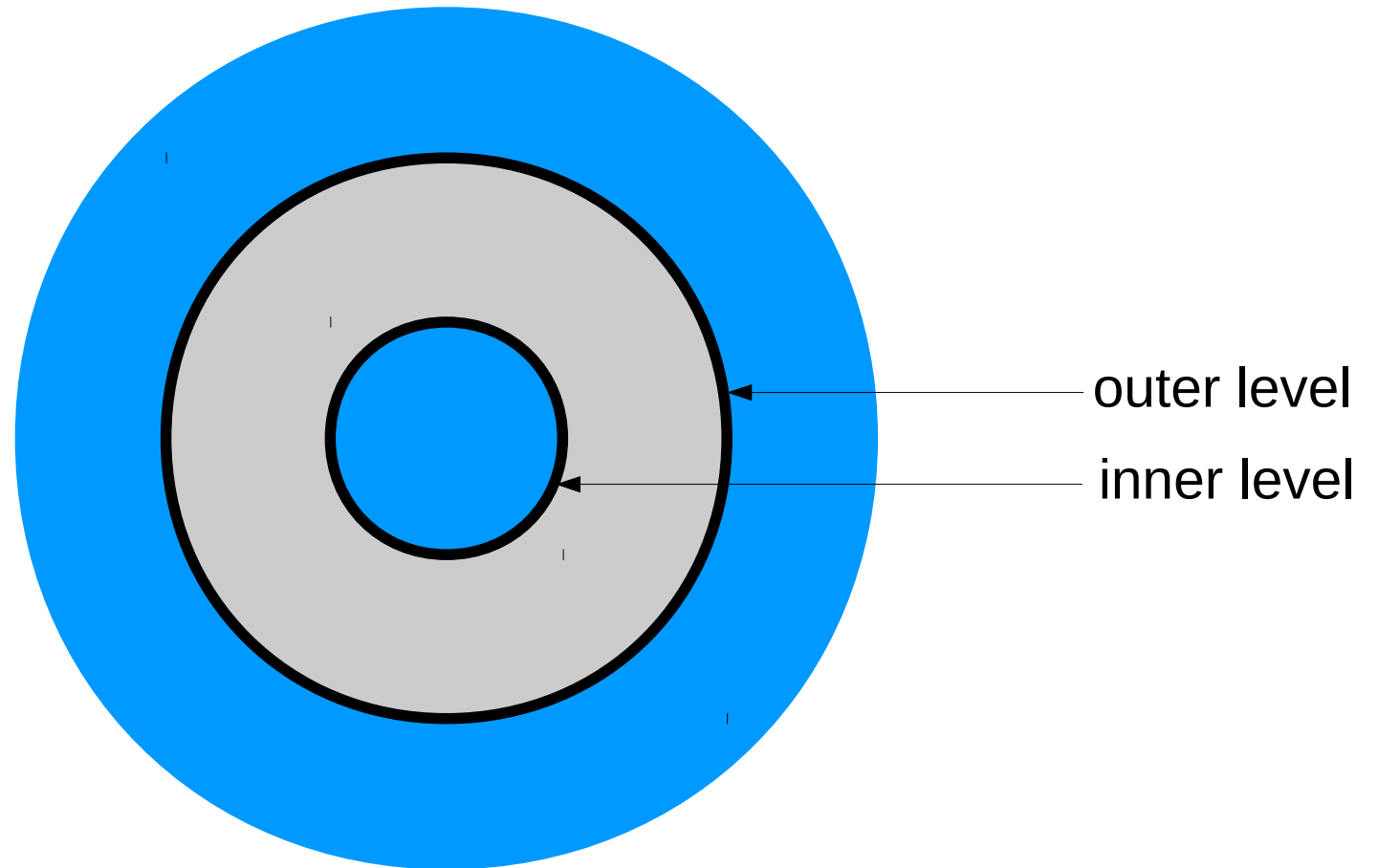




Levels



Levels

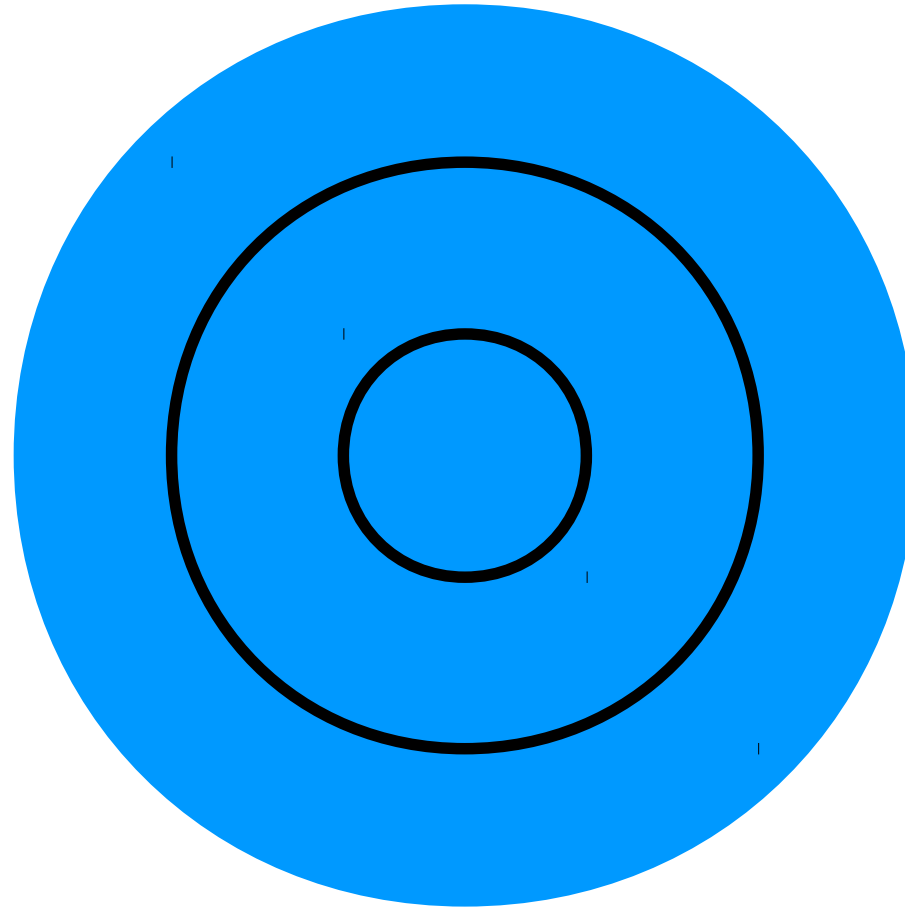


weight $< 2/3$

weight $\geq 2/3$

weight unknown

First Case:

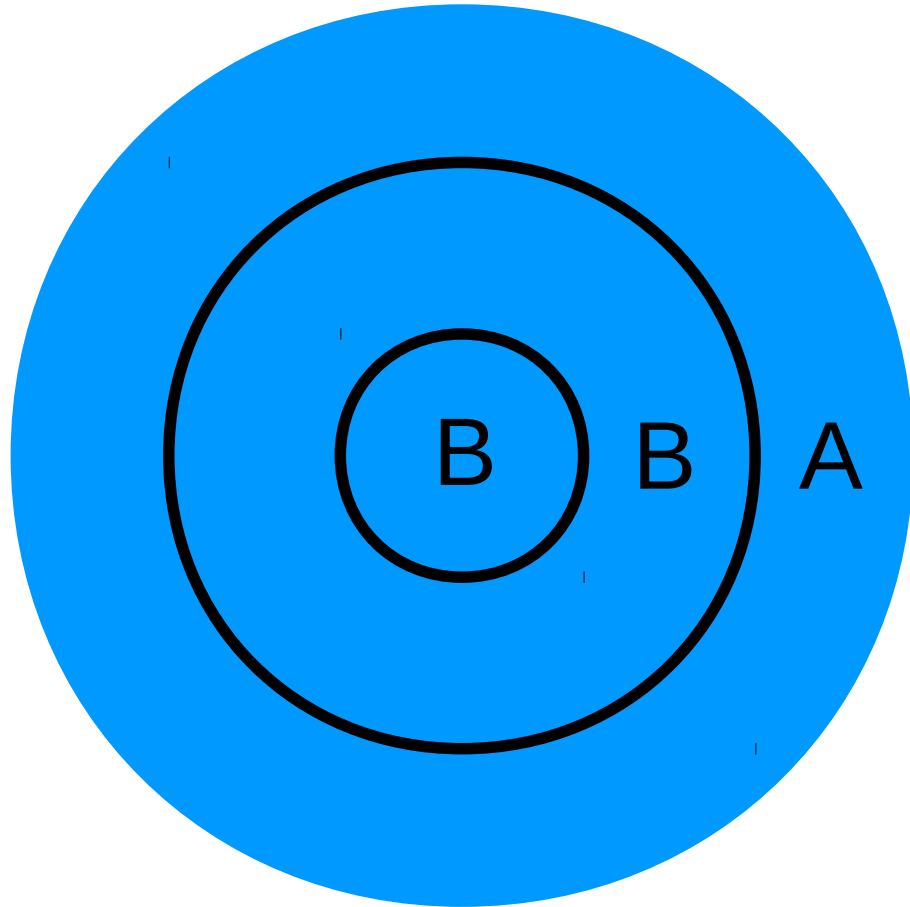


weight $< 2/3$

weight $\geq 2/3$

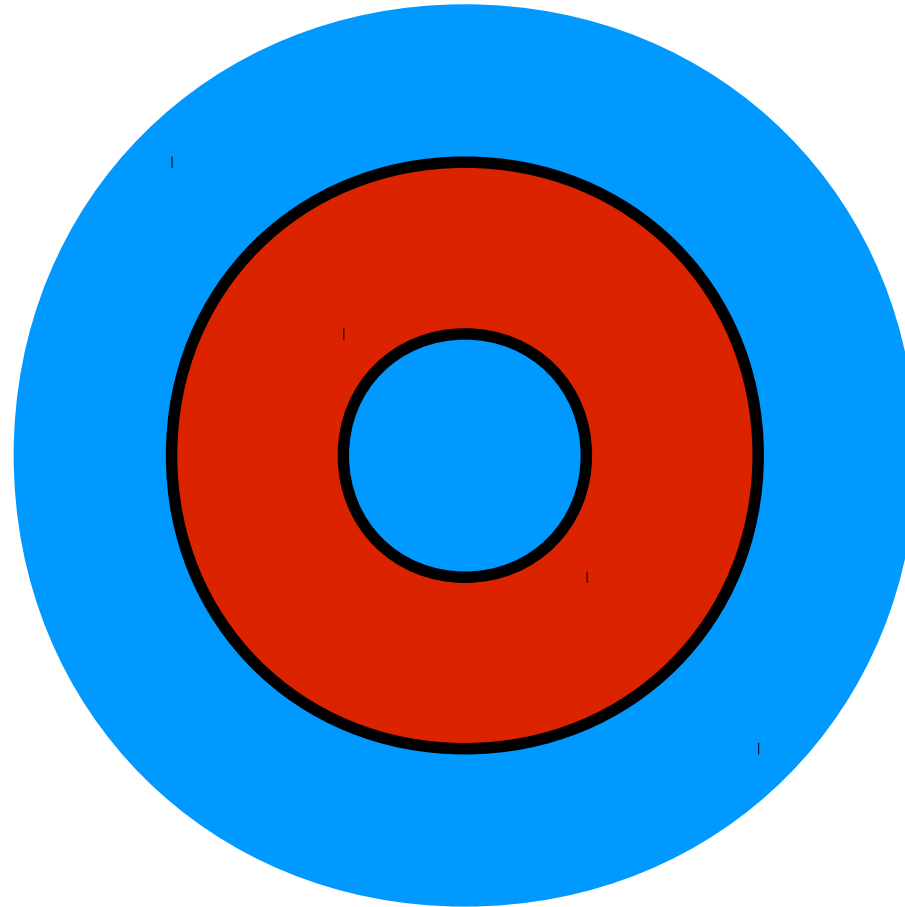
weight unknown

First Case:



- **Choose:**
 - **Set A:** heavier set
 - **Set B:** two lighter sets
 - **Separator:** nodes on inner and outer level
- Weight of A and B is smaller than $\frac{2}{3}$

Second Case:

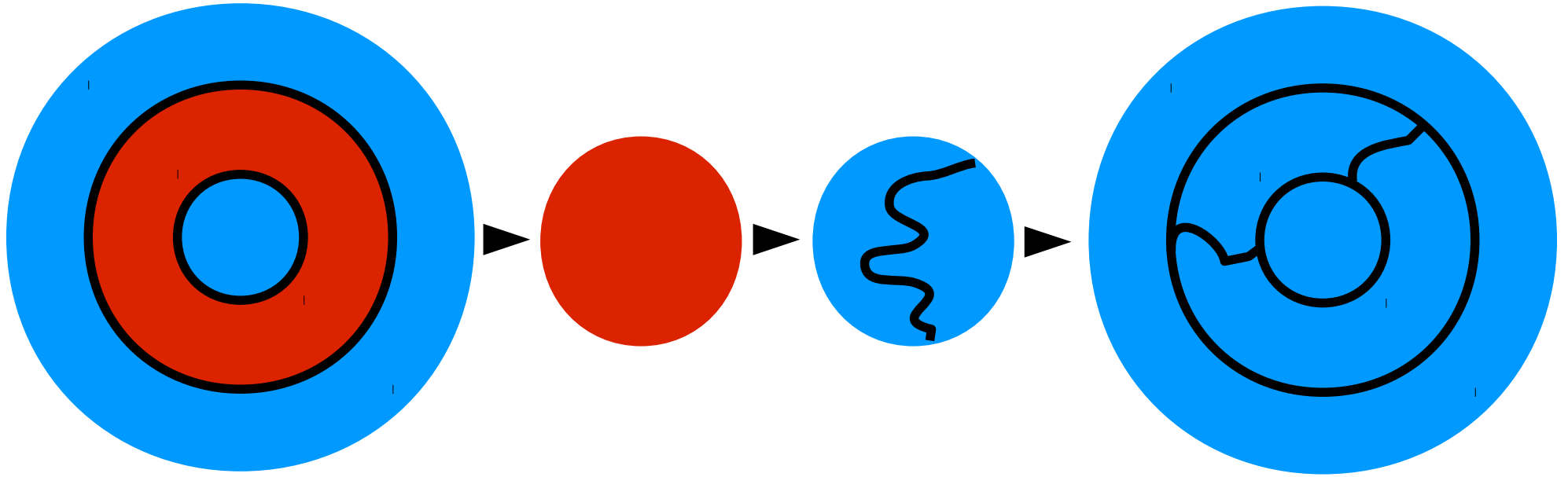


weight $< 2/3$

weight $\geq 2/3$

weight unknown

Second Case:

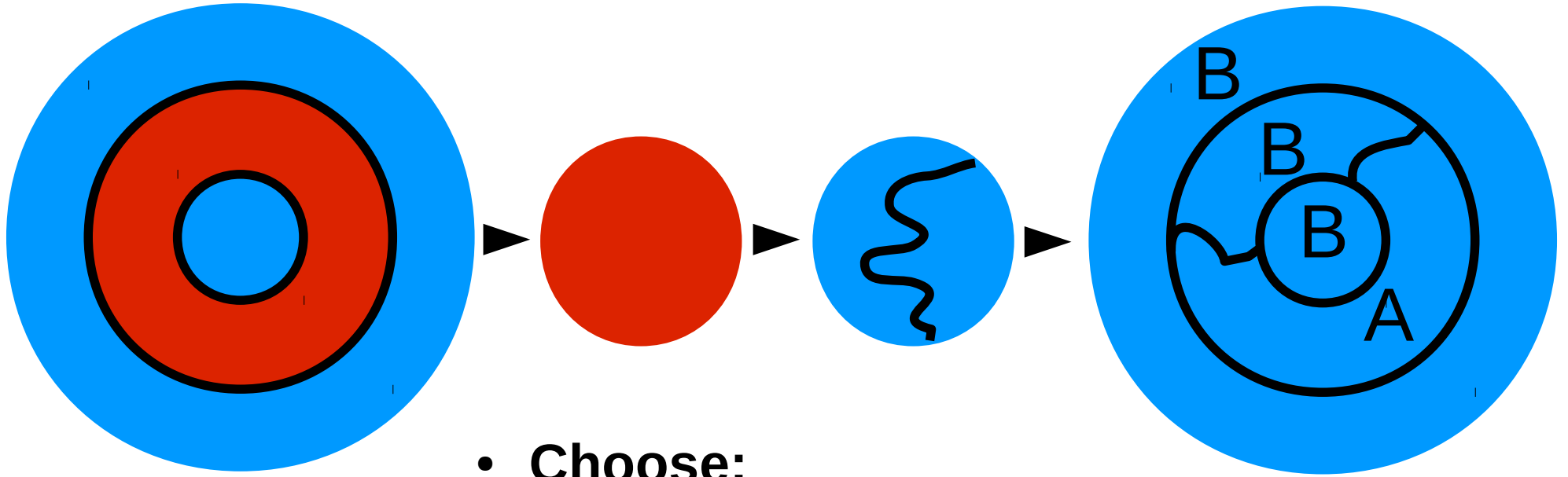


weight $< 2/3$

weight $\geq 2/3$

weight unknown

Second Case:



- **Choose:**
 - **Set A:** heavier half of middle set
 - **Set B:** remaining sets
 - **Separator:** nodes on inner and outer level and separator of middle set

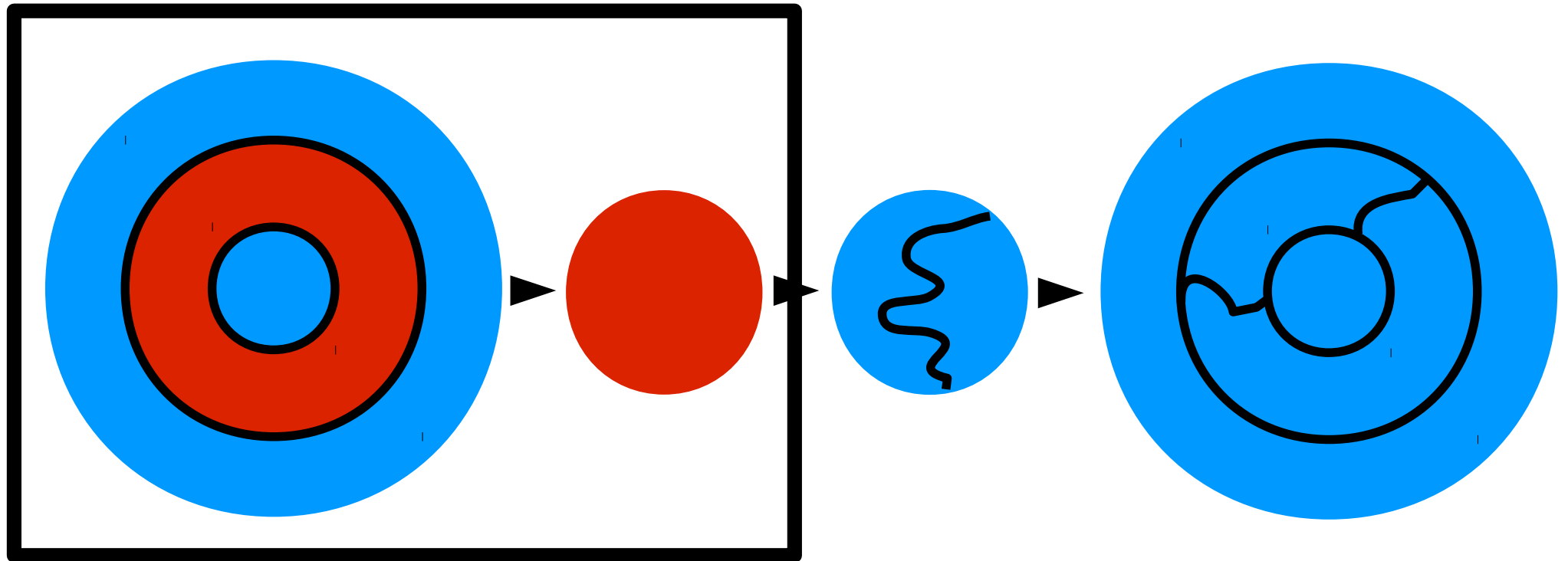
weight $< 2/3$

weight $\geq 2/3$

weight unknown

- Weight of A and B is smaller than $\frac{2}{3}$

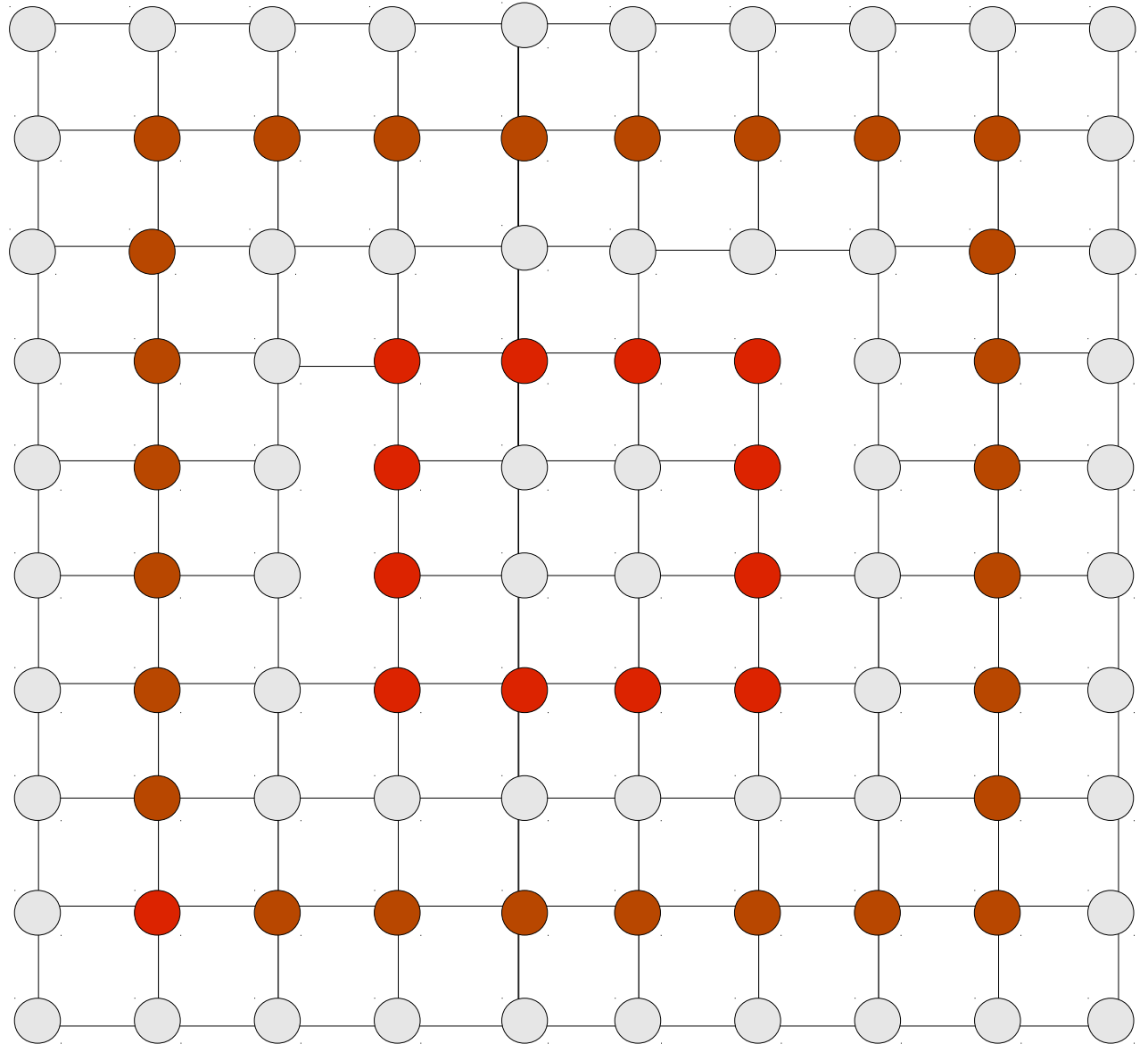
- **How can we separate the middle set?**
 - Extract new graph
 - Find separator
 - Integrate in old graph



- **Idea**

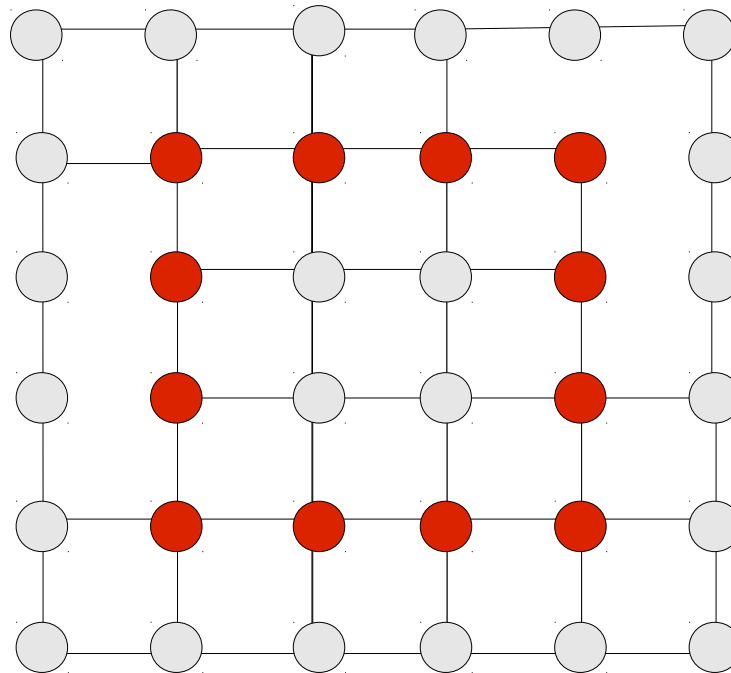
- Remove outer set and separator

- Merge inner set and separator



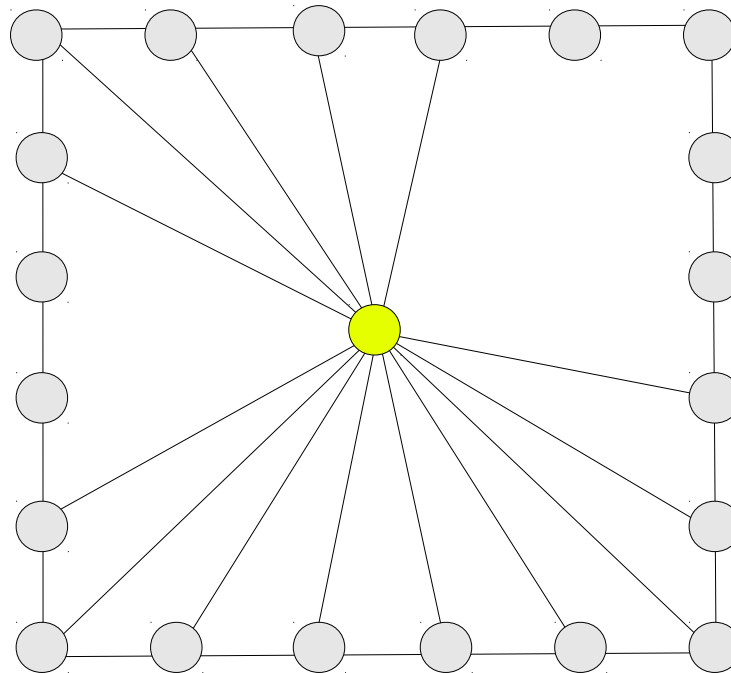
- **Idea**

- Remove outer set and separator
- Merge inner set and separator

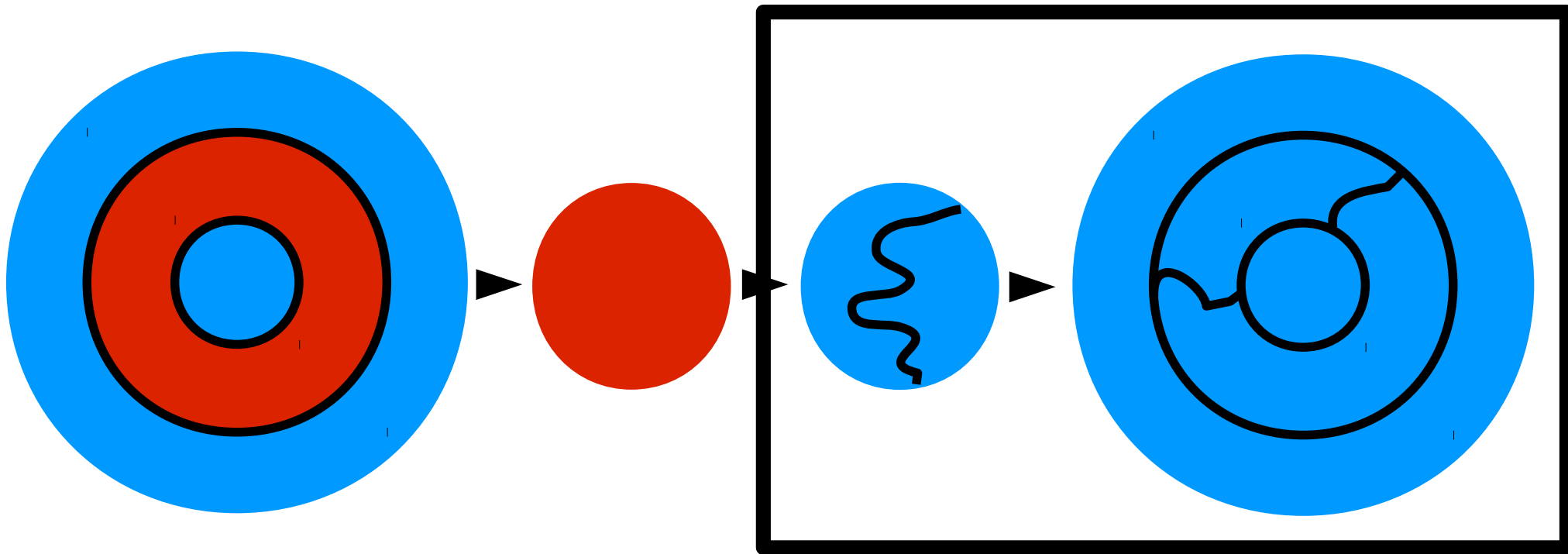


- **Idea**

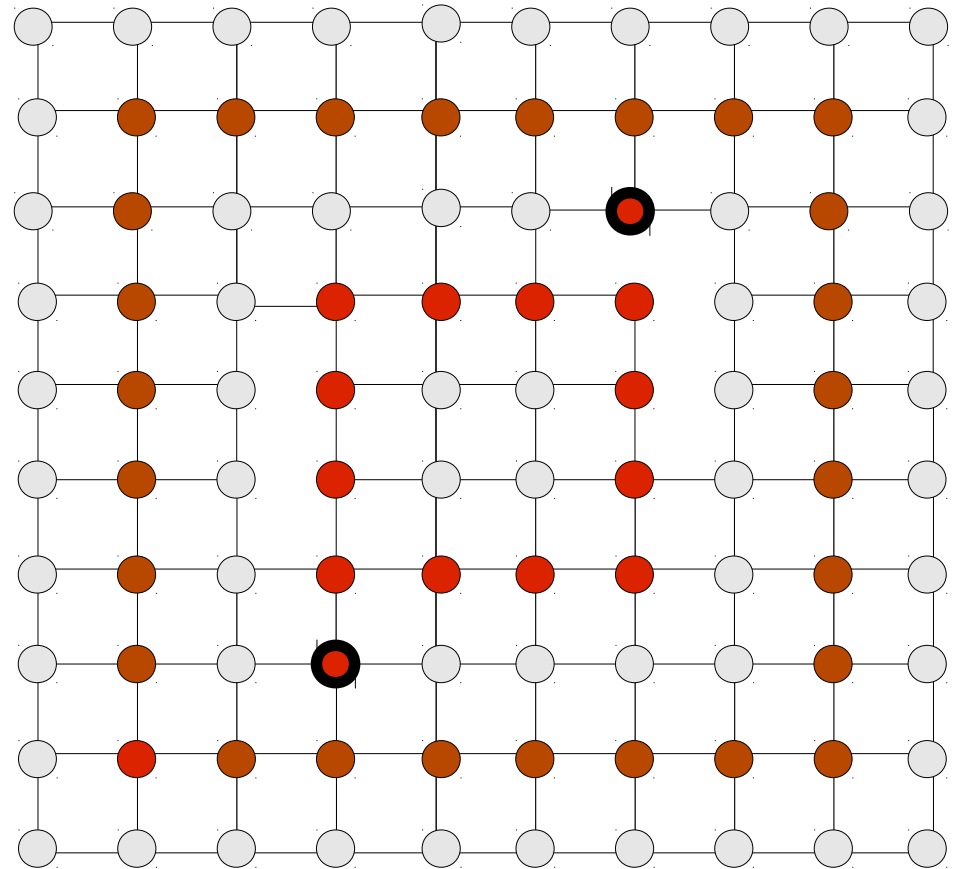
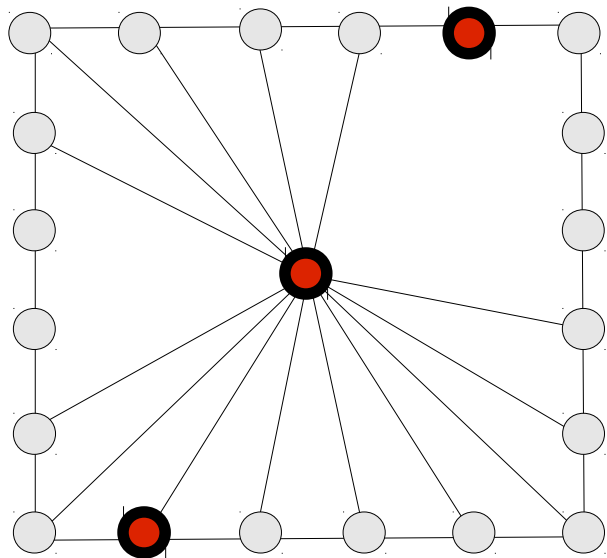
- Remove outer set and separator
- Merge inner set and separator



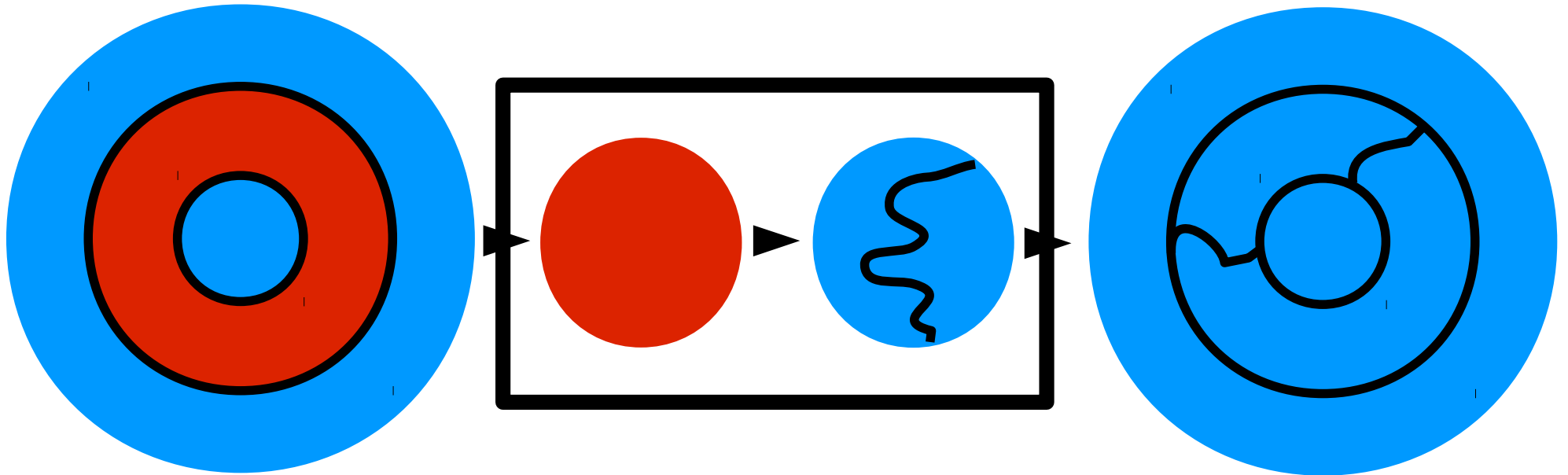
- **How can we separate the middle set?**
 - Extract new graph
 - Find separator
 - Integrate in old graph



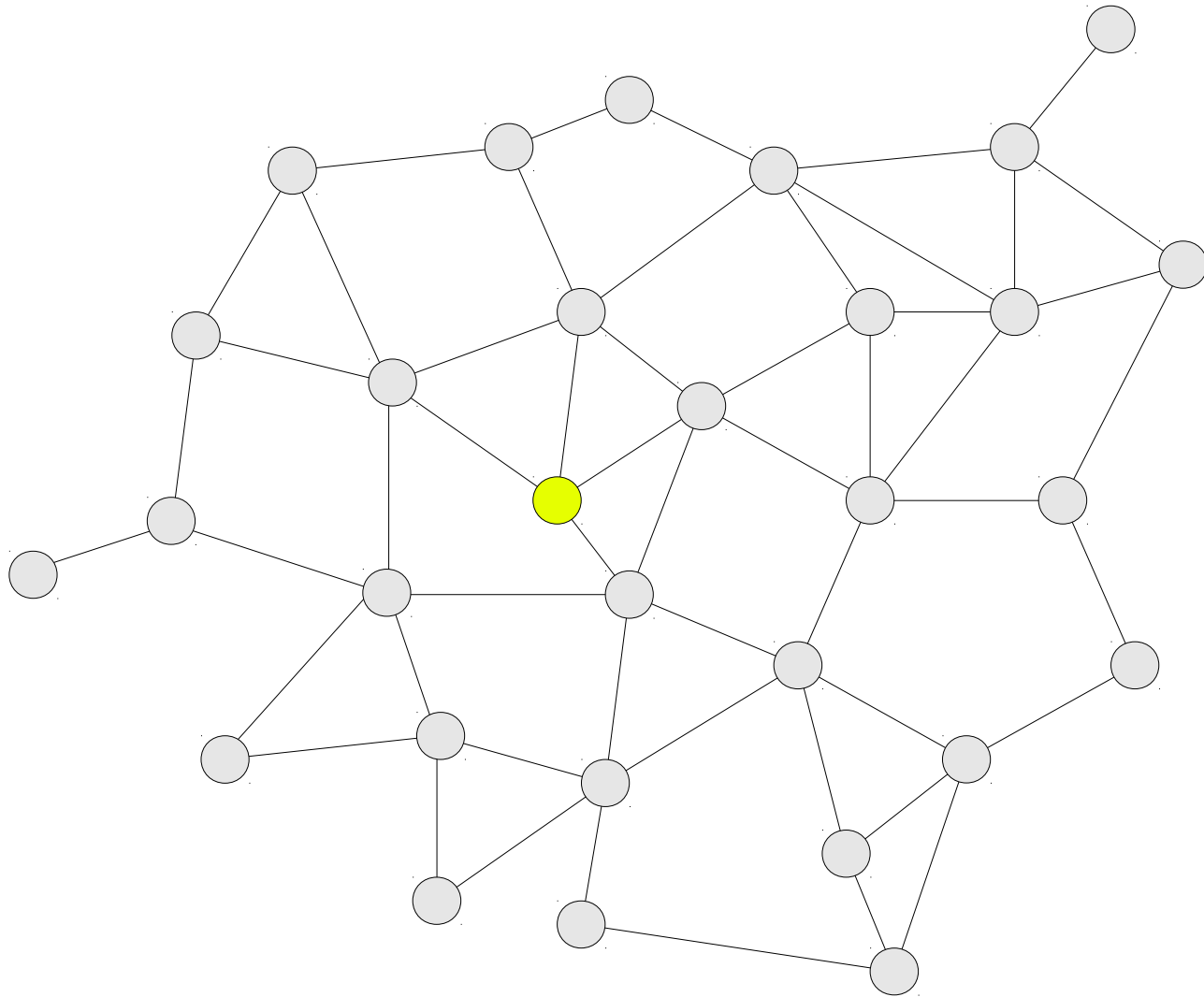
Separating the middle set via spanning tree



- **How can we separate the middle set?**
 - Extract new graph
 - Find separator
 - Integrate in old graph

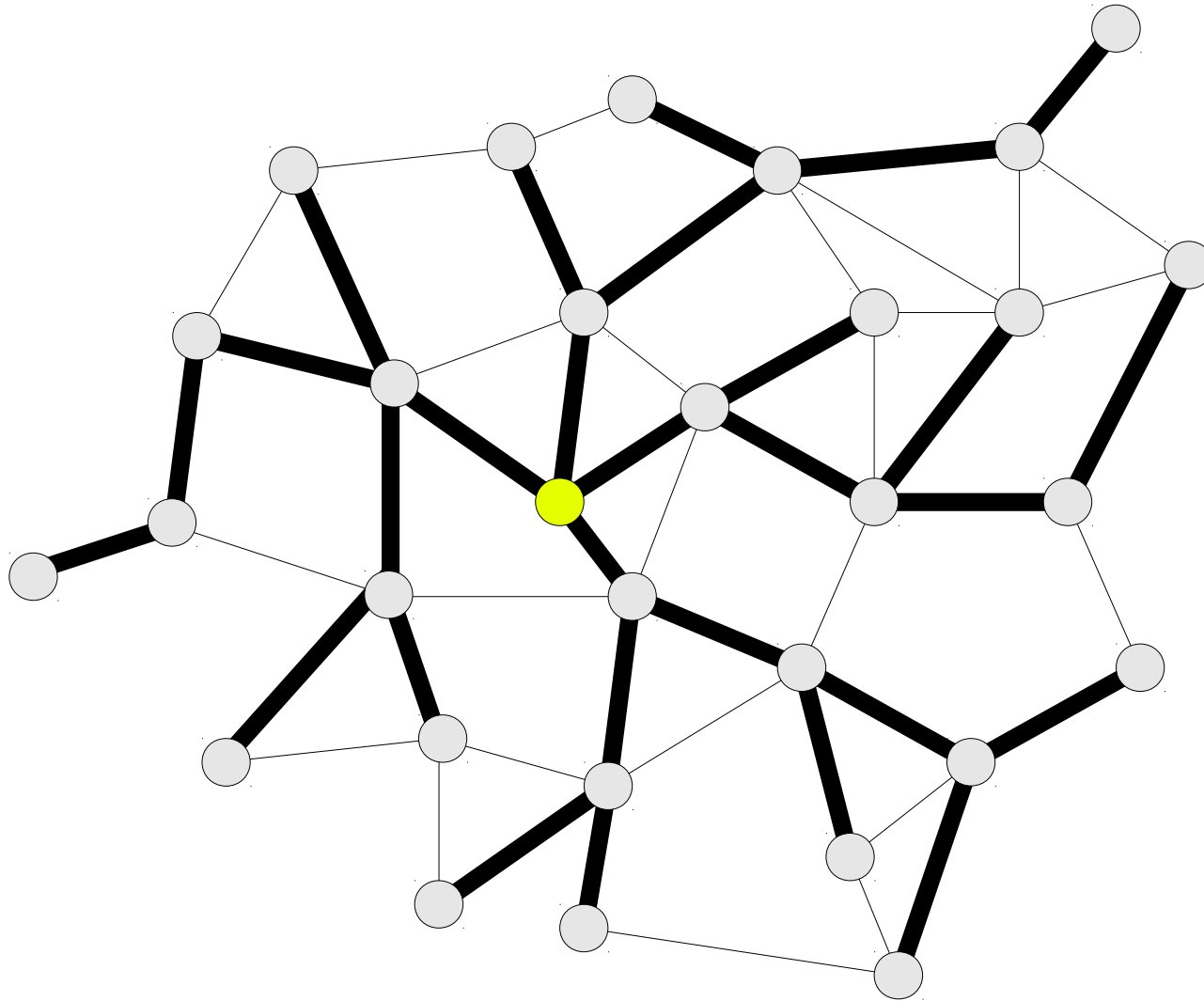


Separating the middle set via spanning tree



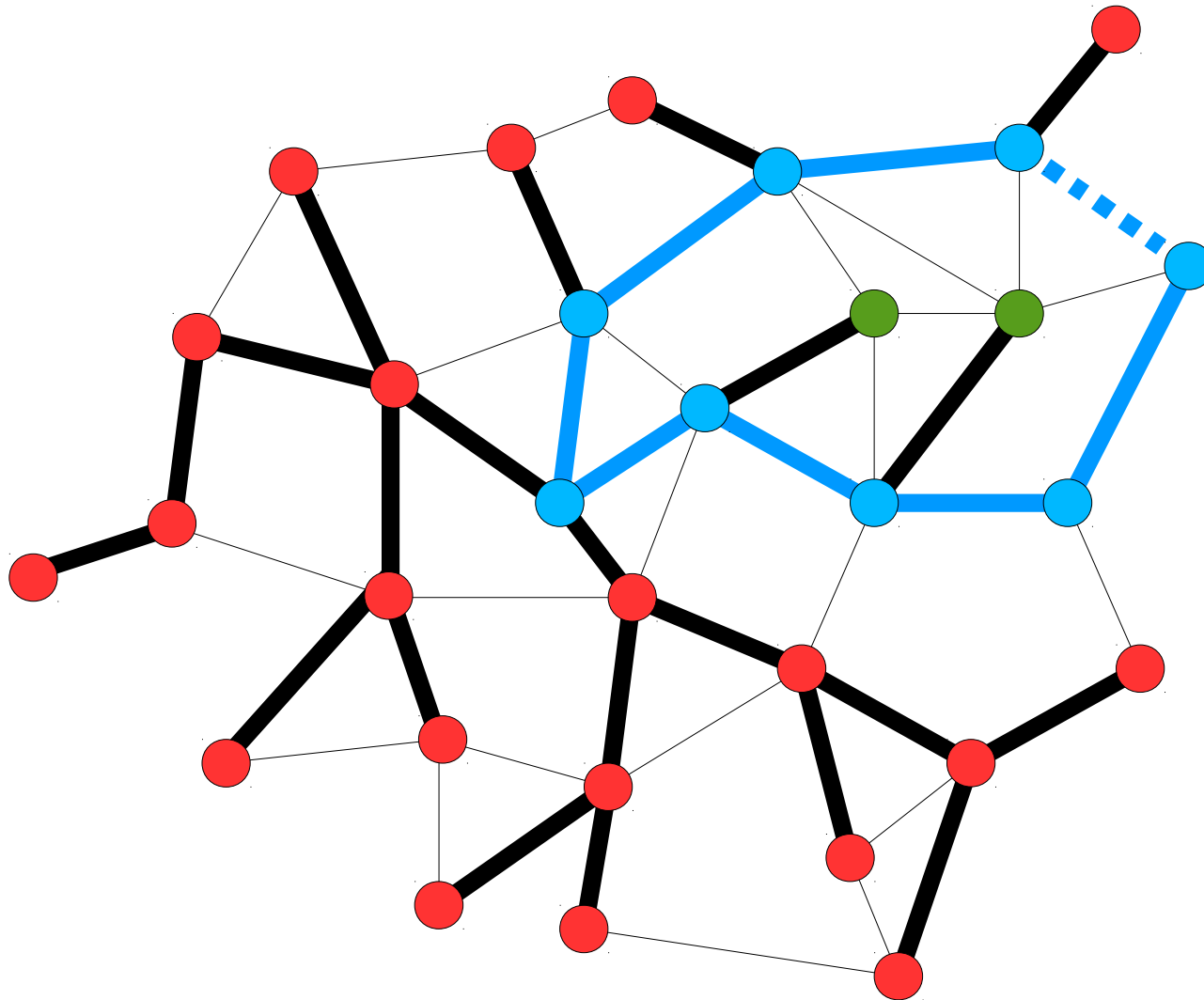
Separating the middle set via spanning tree

- **Calculate:** Planar representation and DFS spanning tree



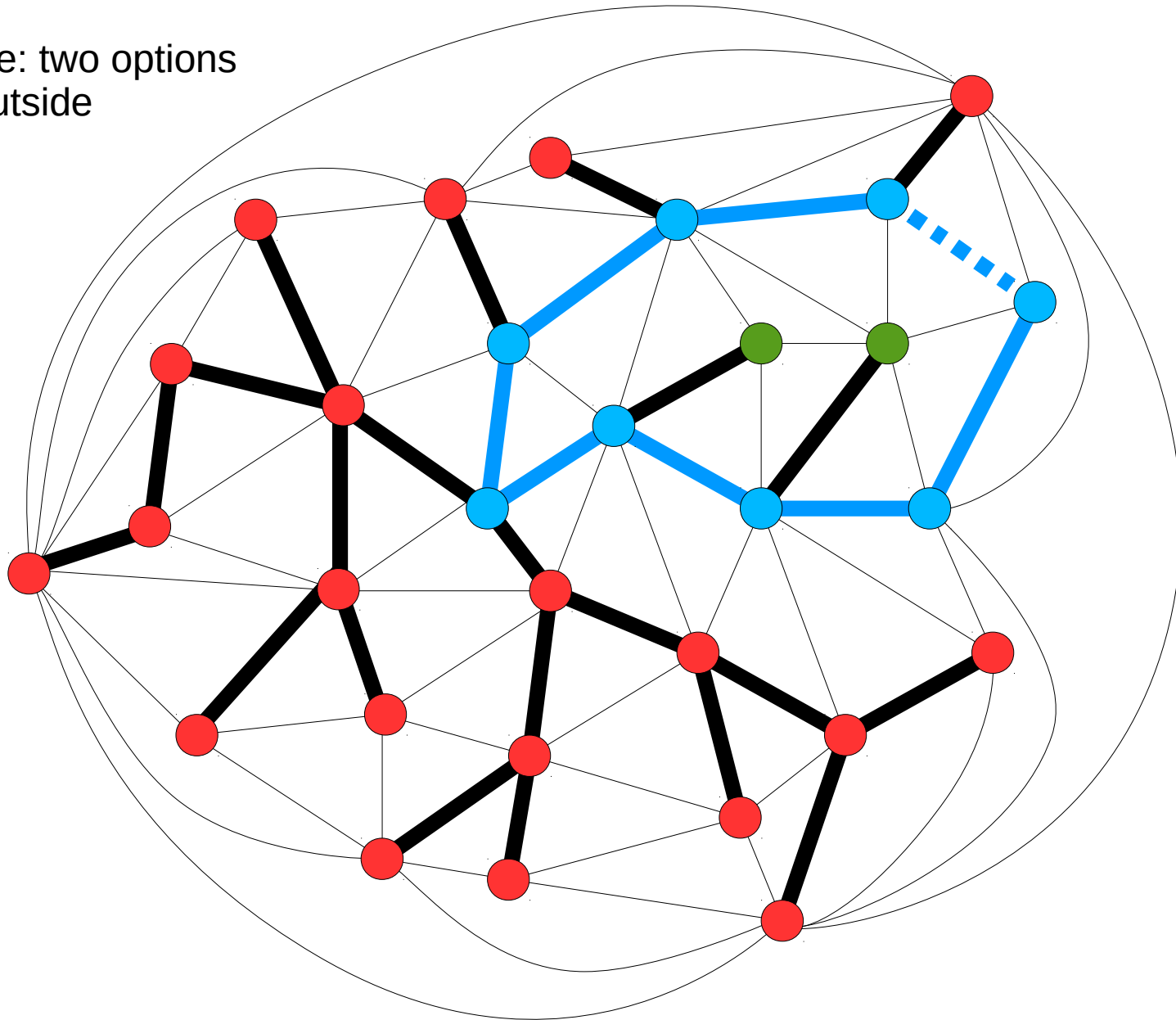
Separating the middle set via spanning tree

- Pick random edge to create a circle

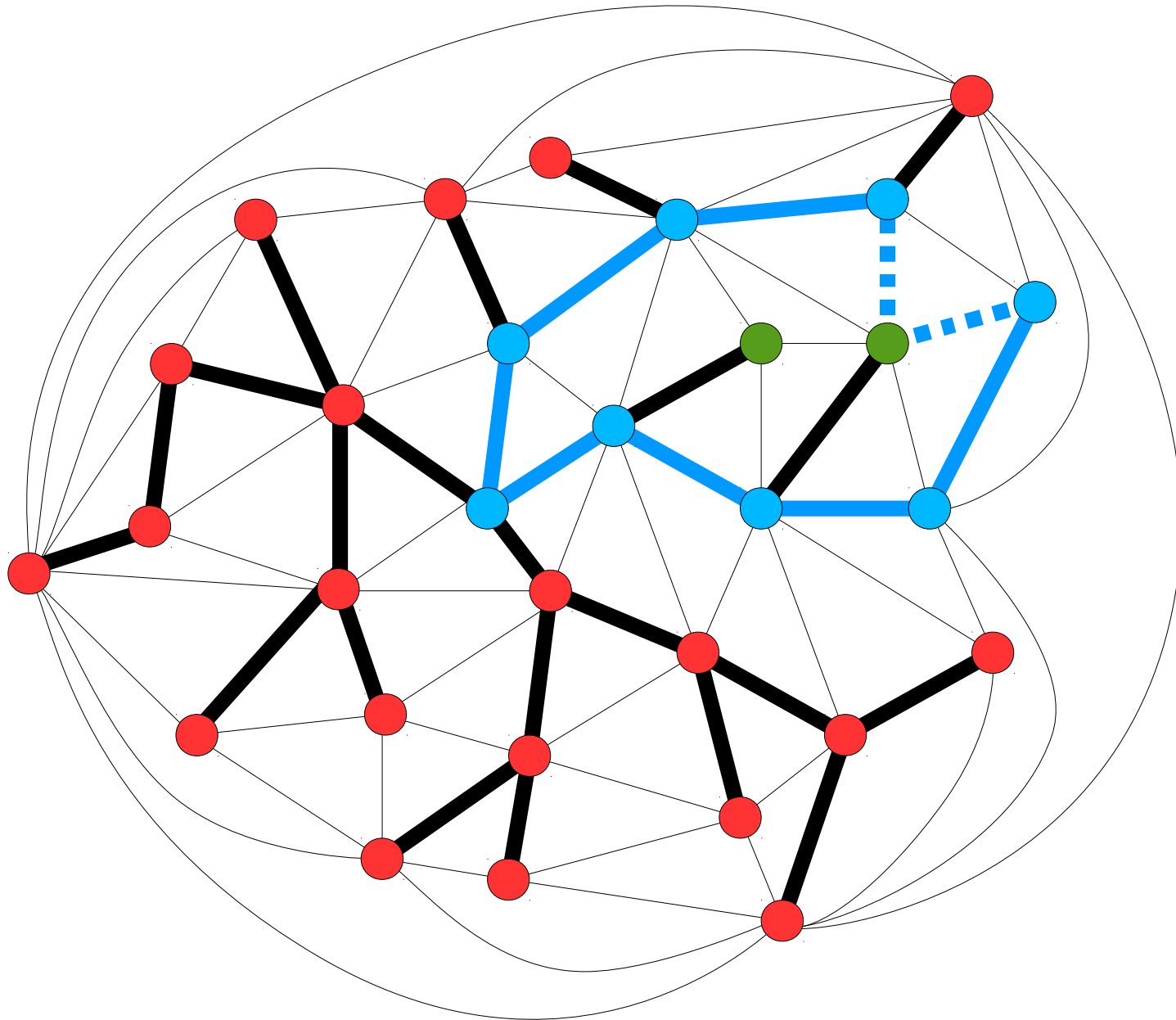


Separating the middle set via spanning tree

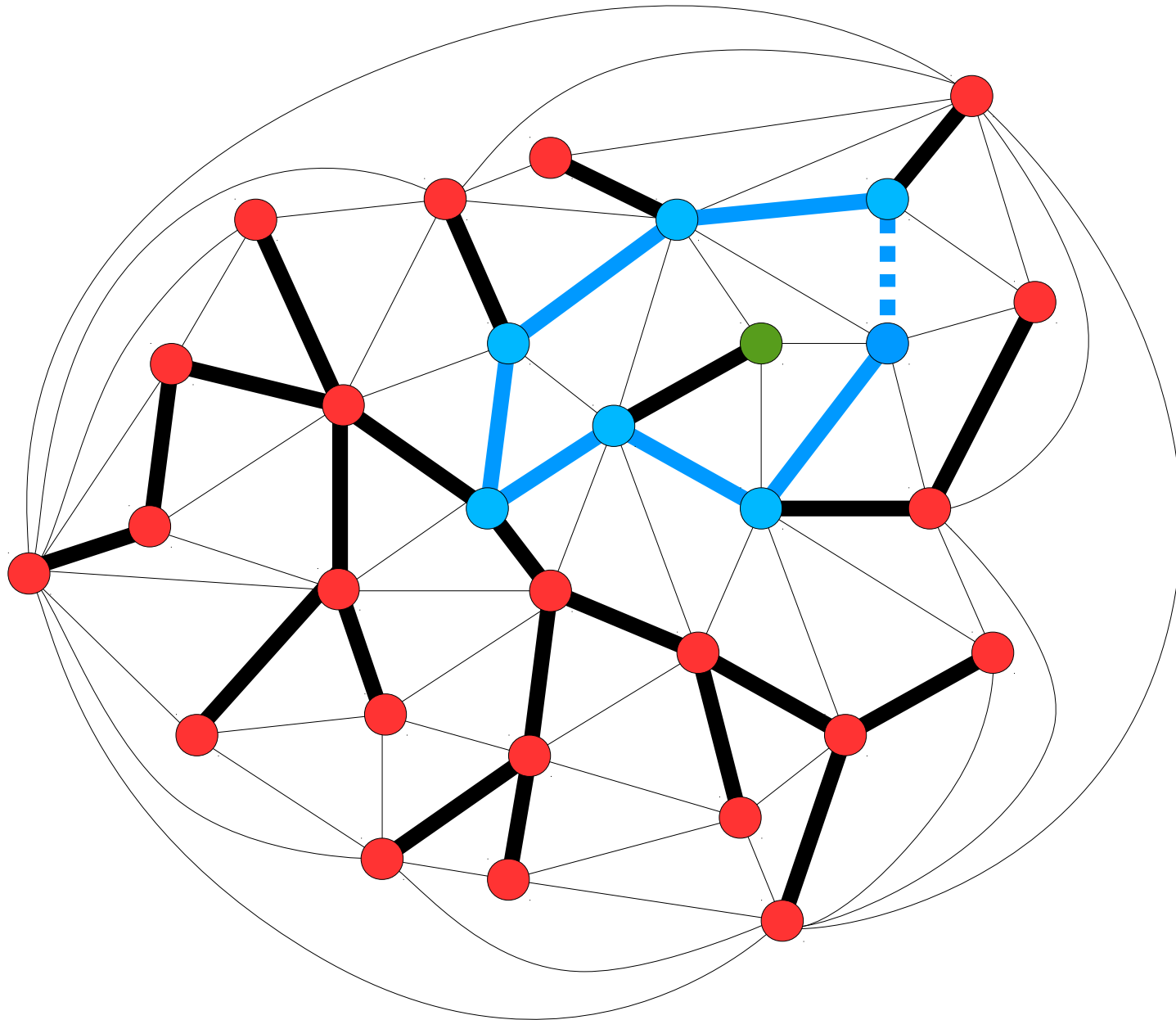
Triangle: two options
Why outside



Separating the middle set via spanning tree

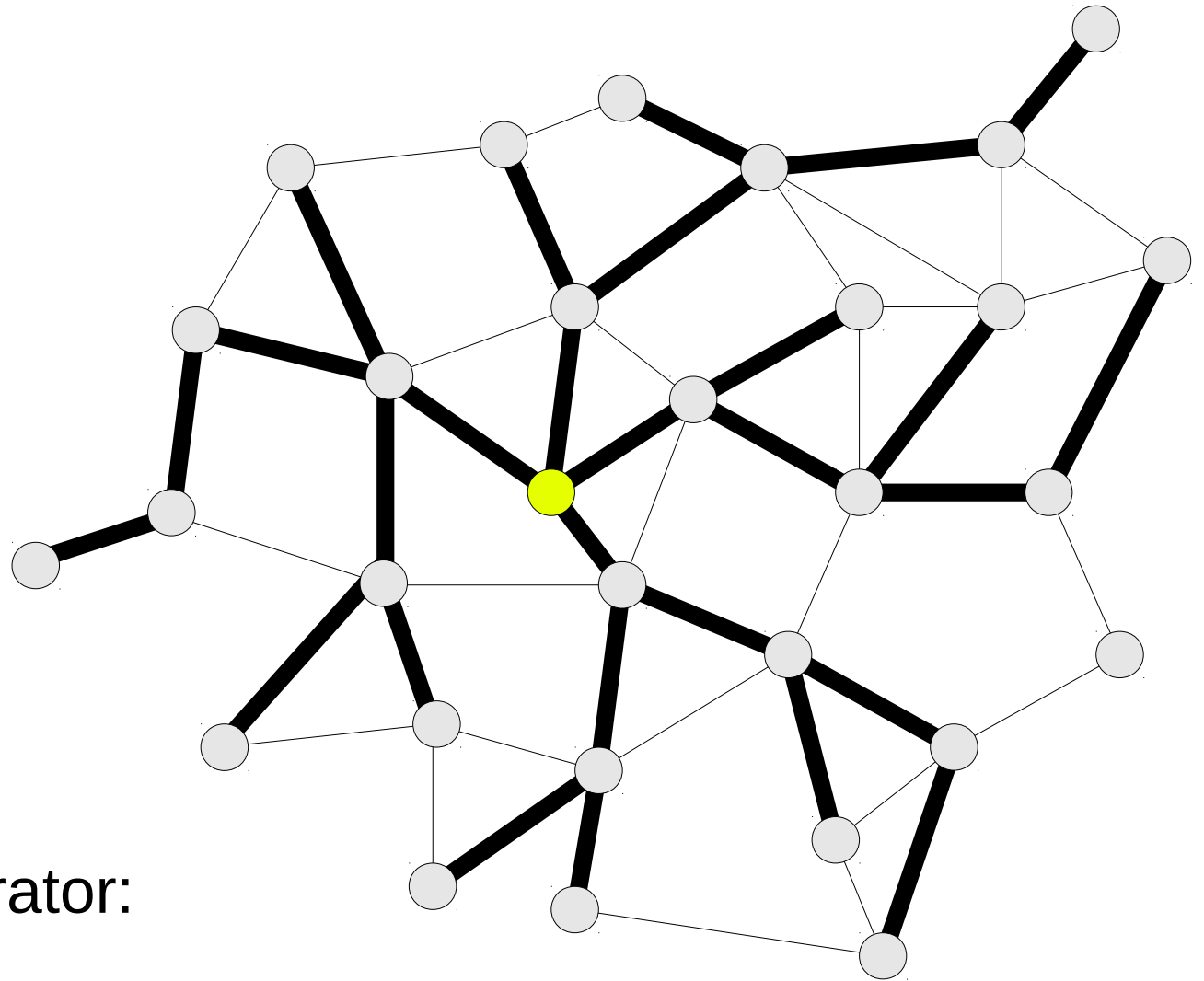
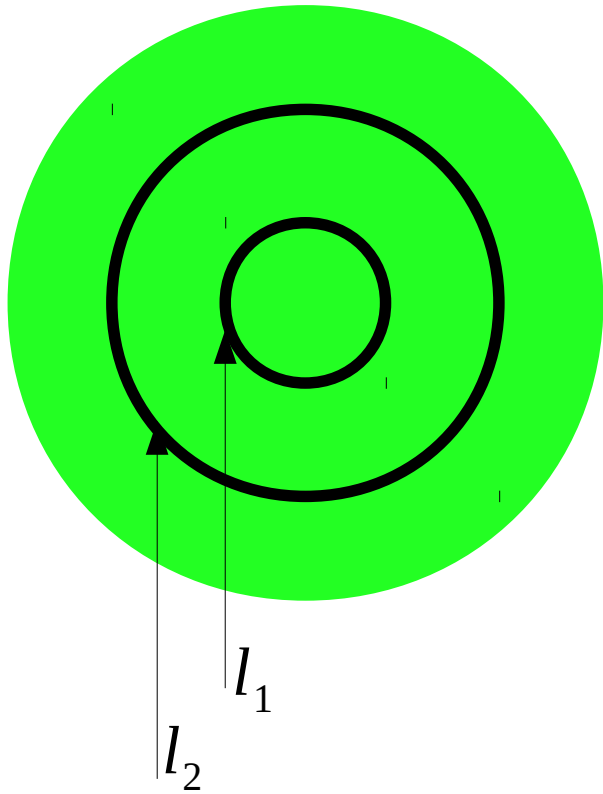


Separating the middle set via spanning tree



- **What is the size of the separator?**

Size of middle separator



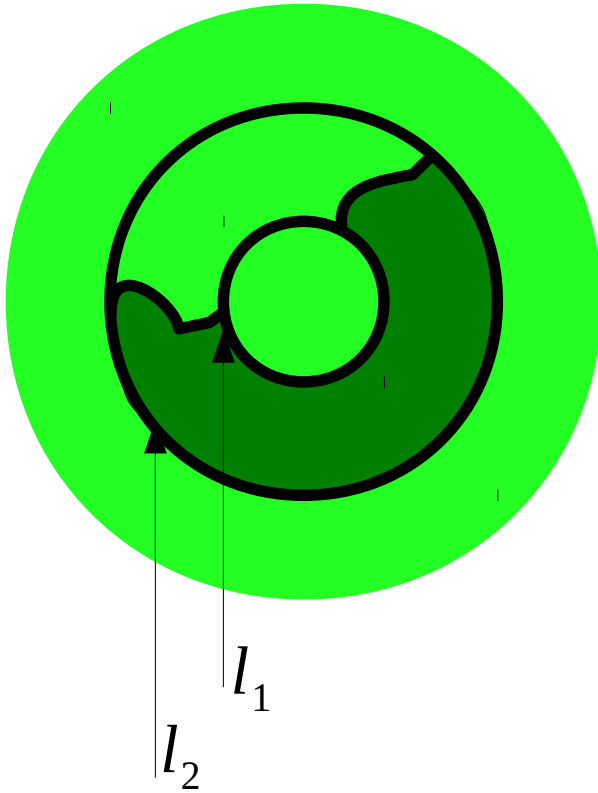
- Radius of tree:

$$r = l_2 - l_1 - 1$$

- Size of middle separator:

$$|S_{middle}| \leq 2(l_2 - l_1 - 1)$$

Size of the whole separator



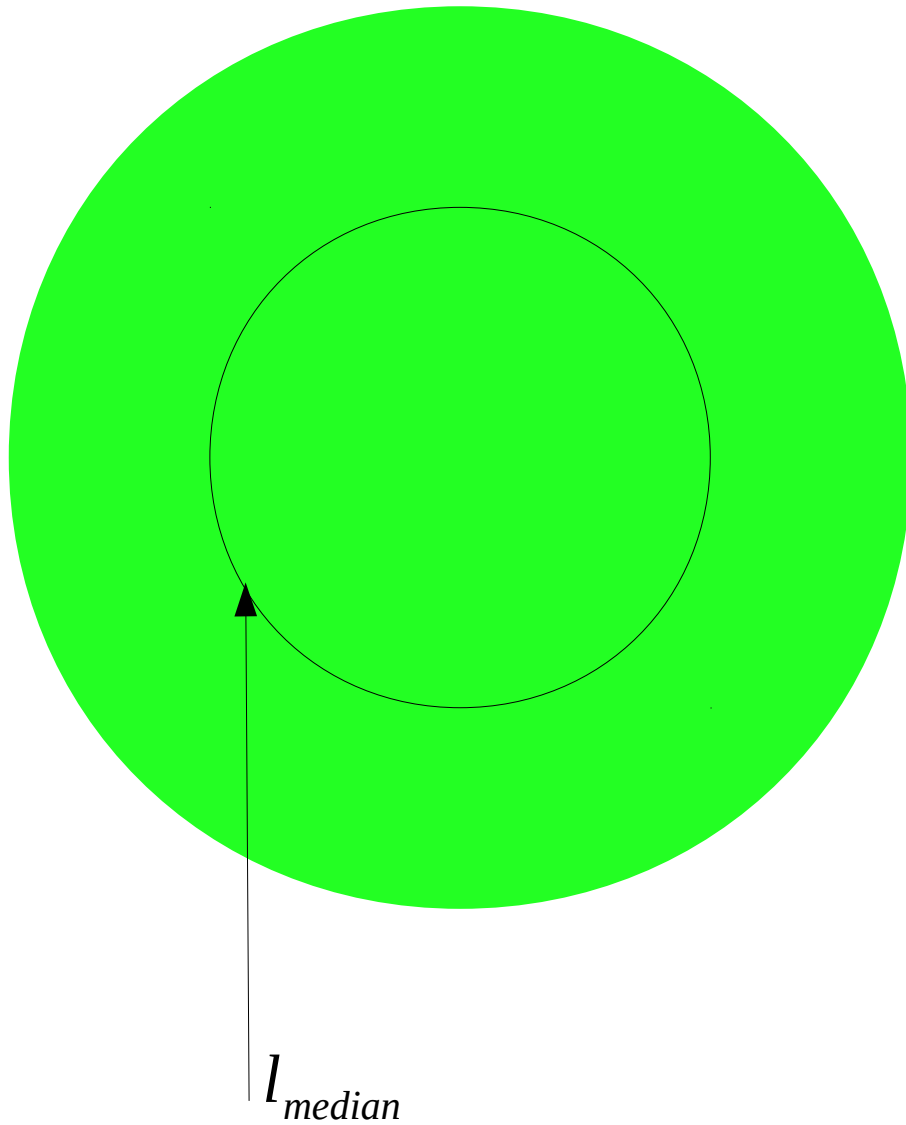
- Separator is linear in number of nodes on both levels plus the middle separator:

$$|S| \leq N(l_1) + N(l_2) + |S_{middle}|$$

$$|S| \leq N(l_1) + N(l_2) + 2(l_2 - l_1 - 1)$$

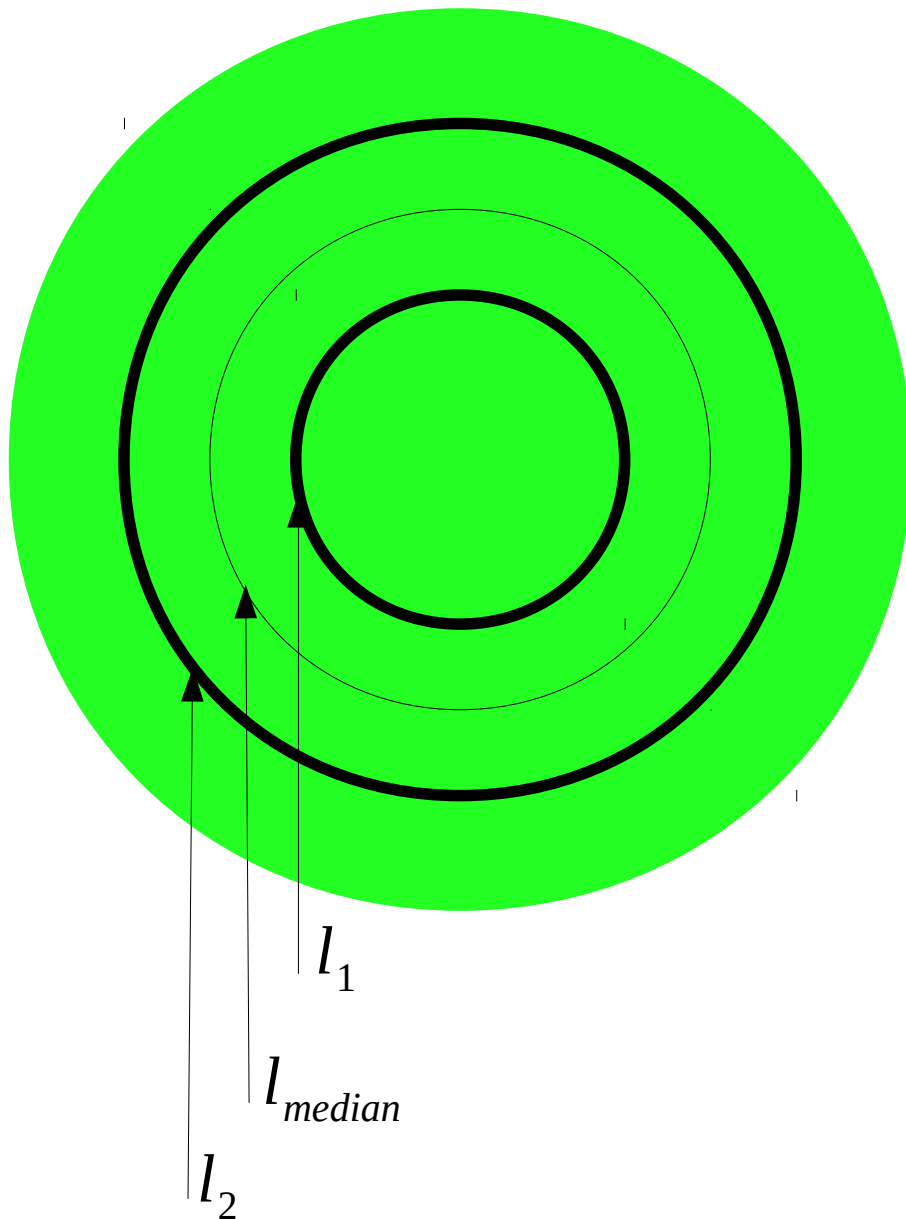
- How to choose l_1 and l_2 so that size of separator is in $O(\sqrt{n})$?

Size of the whole separator



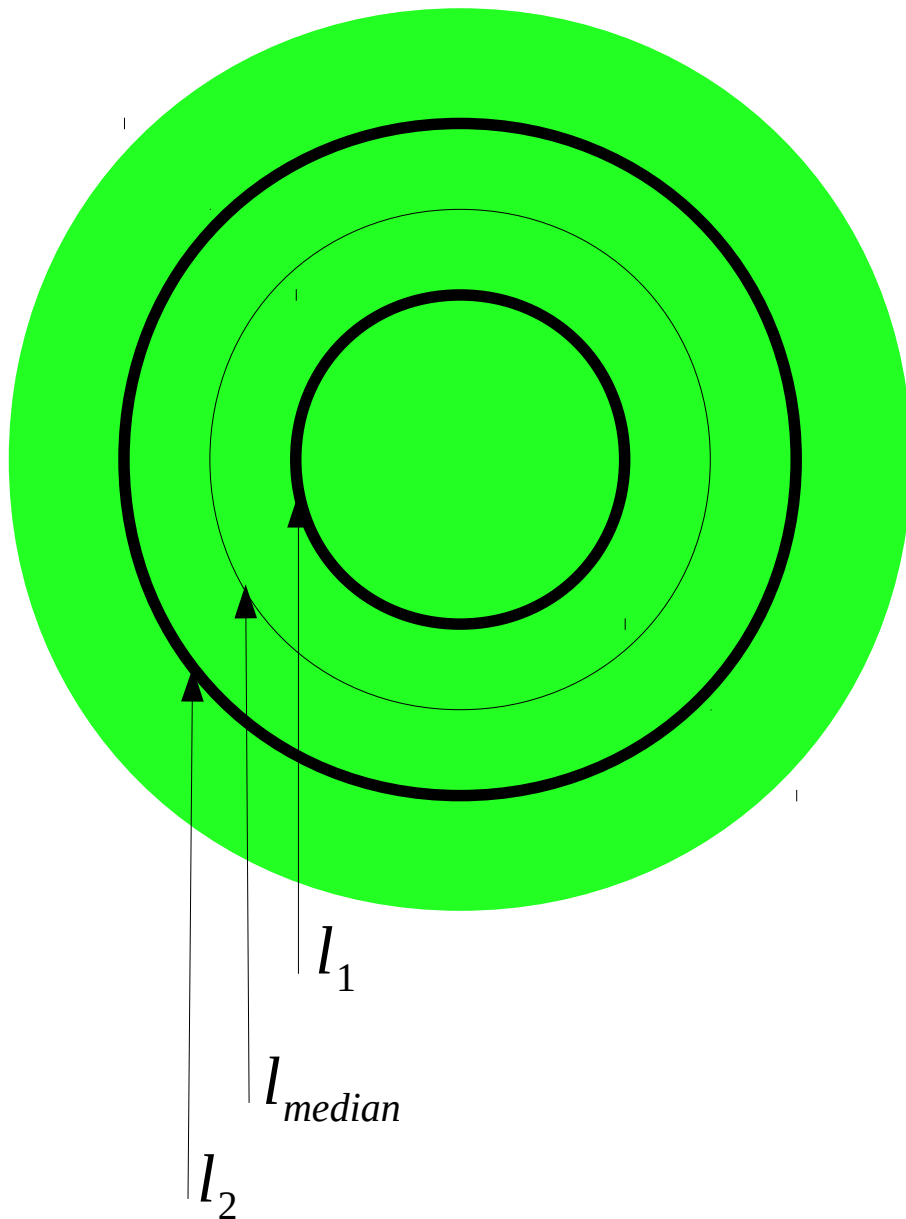
- Find median level l_{median} , so that inner and outer set have weight of less than $\frac{1}{2}$
- $k :=$ number of nodes on this level

Size of the whole separator



- Find median level l_{median} , so that inner and outer set have weight of less than $\frac{1}{2}$
- $k :=$ number of nodes on this level

Size of the whole separator



- Find median level l_{median} , so that inner and outer set have weight of less than $\frac{1}{2}$
- $k :=$ number of nodes on this level
- Choose l_1 , with:
$$N(l_1) + 2(l_{median} - l_1) < 2\sqrt{k}$$
- Choose l_2 , with:
$$N(l_2) + 2(l_2 - l_{median} - 1) < 2\sqrt{(n-k)}$$

Size of the whole separator

$$|S| \leq N(l_1) + N(l_2) + 2(l_2 - l_1 - 1)$$

$$\left| \begin{array}{l} N(l_1) + 2(l_{\text{median}} - l_1) < 2\sqrt{k} \\ N(l_2) + 2(l_2 - l_{\text{median}} - 1) < 2\sqrt{n-k} \end{array} \right.$$

$$|S| \leq 2\sqrt{k} + 2\sqrt{n-k}$$

$$|S| \leq 2\sqrt{n/2} + 2\sqrt{n/2}$$

$$|S| \leq 2\sqrt{2}\sqrt{n}$$

Sources:

- Lipton, Richard J.; Tarjan, Robert E. (1979), "A separator theorem for planar graphs", *SIAM Journal on Applied Mathematics* 36 (2): 177–189