

## Tutorial Exact Algorithms

### Exercise T6 (again)

The problem MAXCUT is defined as follows:

- Input: A graph  $G = (V, E)$ .  
Feasible solutions: Any bipartition  $V = V_1 \cup V_2$ .  
Goal: Maximize the number of edges that are *cut* by the bipartition.

Design an algorithm for this problem with a running time of  $O^*(\tau(3, 3)^m)$ , where as usual  $m = |E|$ .

### Exercise T7

Let  $F = \{C_1, \dots, C_m\}$  be a formula in 3-CNF with clauses  $C_1, \dots, C_m$ . Define for a variable  $x$  the “maxdegree” of  $x$  as  $d(x) := \max\{|C| \mid x \in C\}$

Now consider the following simple branching algorithm for the 3SAT problem:

Given is a formula  $F$  in 3-CNF. If there is  $x$  that occurs only positively or negatively, or a clause  $\{x\}$  or  $\{\bar{x}\}$ , choose the value for  $x$  accordingly. Otherwise choose  $x$  with maximal  $d(x)$  and branch on the two possible assignments  $x := 0$  and  $x := 1$ .

What is the running time in the form  $O^*(c^n)$  when analyzed in the number  $n$  of variables?

Now consider a new measure. Let  $\alpha_1, \alpha_2, \alpha_3 \in \mathbf{R}$  and define

$$\Phi(F) := \sum_{x \in \text{vars}(F)} \alpha_{d(x)}.$$

- Construct the branching vectors for  $\Phi(F)$  for all possible cases.
- What are good choices for  $\alpha_1, \alpha_2, \alpha_3$ ?
- What running time do we obtain?

### Homework Assignment H7 (10 Points)

Consider a formula in  $F$  in 4-CNF and the same algorithm as in Exercise T6. Let again  $\alpha_1, \alpha_2, \alpha_3, \alpha_4 \in \mathbf{R}$  and

$$\Phi(F) := \sum_{x \in \text{vars}(F)} \alpha_{\delta(x)}.$$

Compute, in terms of  $\alpha_1, \alpha_2, \alpha_3$  and  $\alpha_4$ , the branching vectors for all possible cases when measured in  $\Phi(F)$ .

### Homework Assignment H8 (10 Points)

Find optimal values for  $\alpha_1, \alpha_2, \alpha_3$  and  $\alpha_4$ .