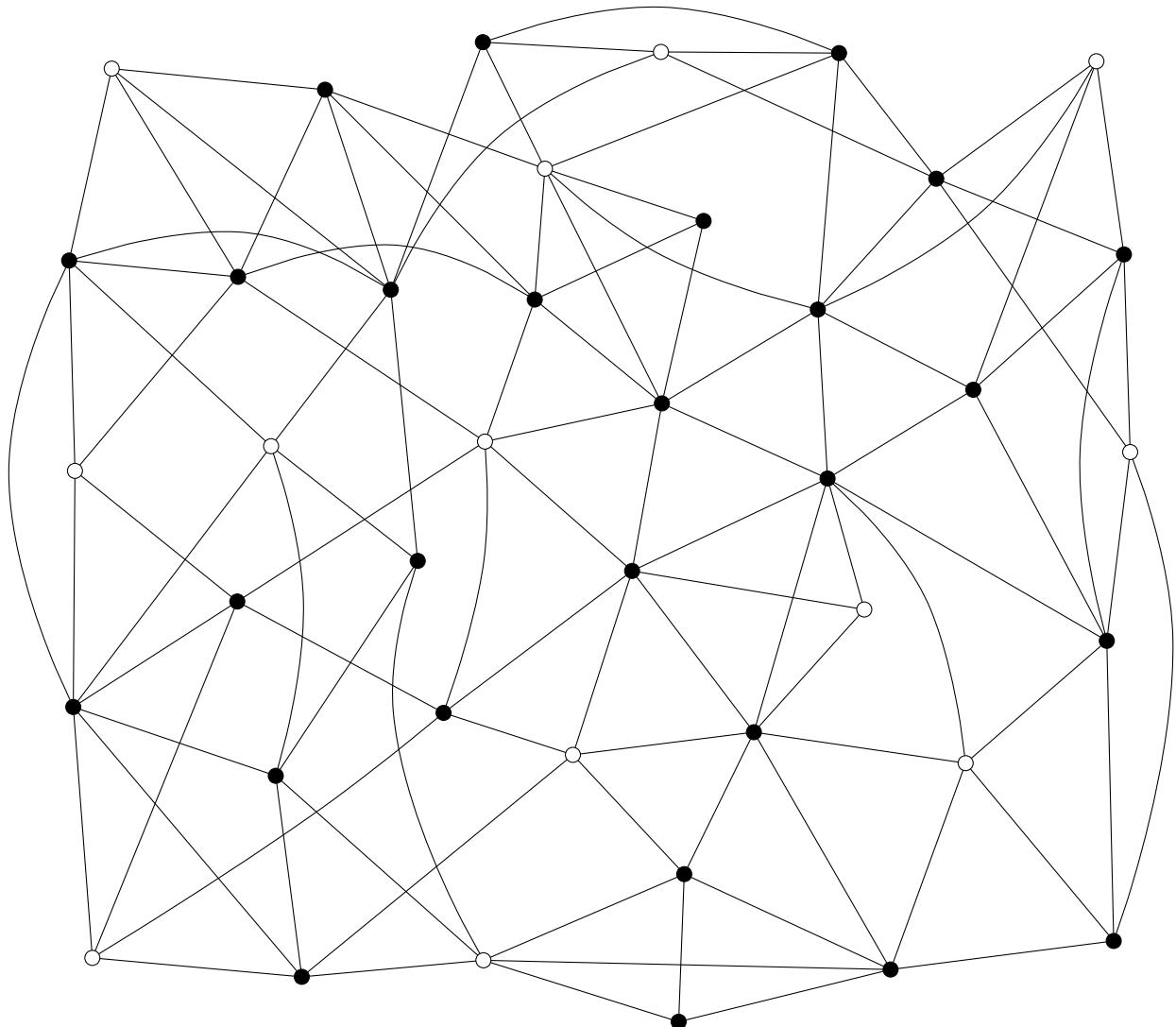


Exercise Sheet 03

Problem T7

A *vertex cover* of an undirected graph $G = (V, E)$ is a subset $C \subseteq V$ of its vertices such that at least one endpoint of every edge is in C , i.e., for every $(v_i, v_j) \in E$, either $v_i \in C$ or $v_j \in C$. Informally speaking, “ C covers all edges.” It is an NP-complete problem to find out whether a graph has a vertex cover of a given size. The following example shows a graph of order 40. The black vertices comprise a minimum size vertex cover.



Now, given a graph $G = (V, E)$ we define its size as the number of vertices $|V|$ of the graph. So, let us consider a graph of size n and let k be targeted vertex cover size. Find an algorithm for Vertex Cover that runs in time $1.5^k n^{O(1)}$. Hints: If a graph has maximum degree two, i.e., without any vertex of degree three or more, this problem can be solved in polynomial time. On the other hand, if a vertex v is not in the vertex cover, all of its neighbors have to be there.

Problem T8

Solve the recurrence relation

$$a_n = 5a_{n-1} - 8a_{n-2} + 4a_{n-3}$$

with the starting conditions $a_0 = 1$, $a_1 = 3$, $a_2 = 8$.

Problem H6 (10 credits)

Solve the following recurrence: Let $a_0 = 1$, $a_1 = 1$, $a_2 = 4$ and

$$a_n = 2a_{n-1} - a_{n-2} + 2a_{n-3}, \text{ for } n \geq 3.$$

Problem H7 (10 credits)

Our task is to generate a word of length n over the alphabet $\{0, 1\}$ which contains neither two consecutive zeros nor three consecutive ones.

Daniel proposes the following algorithm: The algorithm generates a word of length n uniformly at random. If the word fulfills the property, it is returned. Otherwise, the algorithm tries again until it finds one.

What is the expected number of rounds the algorithm needs? Consider this in particular for 32bit words.