

### Exercise for Analysis of Algorithms

#### Exercise T20

In this exercise we consider the following (regular) CFG  $G$ :

$$\begin{aligned} S &\rightarrow abA \mid bS \mid a \\ A &\rightarrow bA \mid aS \end{aligned}$$

1. Find a generating function for number of words  $s_n$  in  $L(G)$  that have length  $n$ .
2. What is the dominant singularity and what kind of singularity is it?
3. What is the exponential growth of  $s_n$ ?
4. How precisely can you estimate  $s_n$  with just the knowledge of the dominating singularity and its nature?
5. Find a closed formula for  $s_n$  with an additive error of at most  $O(0.8^n)$ .

#### Solution:

1. Since the grammar is unique, the symbolic method gives us

$$S(z) = z^2A(z) + zS(z) + z,$$

$$A(z) = zA(z) + zS(z).$$

We solve the latter for  $A(z)$  and obtain  $A(z) = zS(z)/(1 - z)$ . Now we can insert it into the former. This yields

$$S(z) = z^3S(z)/(1 - z) + zS(z) + z.$$

We then solve for  $S(z)$  and get the generating function

$$S(z) = \frac{z}{1 - z - z^3/(1 - z)} = \frac{z(1 - z)}{(1 - z)^2 - z^3}.$$

2. The singularities are the roots of the denominator. We ask Maxima

solve $((1-z)^2 - z^3, z)$  and get

$$z = -\frac{(9\sqrt{23} + 11\sqrt{3})^{\frac{2}{3}} (\sqrt{3}i + 1) + 5 \cdot 2^{\frac{2}{3}} \cdot 3^{\frac{5}{6}} i - 2^{\frac{4}{3}} \cdot 3^{\frac{1}{6}} (9\sqrt{23} + 11\sqrt{3})^{\frac{1}{3}} - 5 \cdot 2^{\frac{2}{3}} \cdot 3^{\frac{1}{3}}}{2^{\frac{4}{3}} \cdot 3^{\frac{7}{6}} (9\sqrt{23} + 11\sqrt{3})^{\frac{1}{3}}},$$

$$z = \frac{(9\sqrt{23} + 11\sqrt{3})^{\frac{2}{3}} (\sqrt{3}i - 1) + 5 \cdot 2^{\frac{2}{3}} \cdot 3^{\frac{5}{6}} i + 2^{\frac{4}{3}} \cdot 3^{\frac{1}{6}} (9\sqrt{23} + 11\sqrt{3})^{\frac{1}{3}} + 5 \cdot 2^{\frac{2}{3}} \cdot 3^{\frac{1}{3}}}{2^{\frac{4}{3}} \cdot 3^{\frac{7}{6}} (9\sqrt{23} + 11\sqrt{3})^{\frac{1}{3}}},$$

$$z = \frac{(9\sqrt{23} + 11\sqrt{3})^{\frac{2}{3}} + 2^{\frac{1}{3}} \cdot 3^{\frac{1}{6}} (9\sqrt{23} + 11\sqrt{3})^{\frac{1}{3}} - 5 \cdot 2^{\frac{2}{3}} \cdot 3^{\frac{1}{3}}}{2^{\frac{1}{3}} \cdot 3^{\frac{7}{6}} (9\sqrt{23} + 11\sqrt{3})^{\frac{1}{3}}}$$

Wolfram Alpha even gives us a nice diagram

[http://www.wolframalpha.com/input/?i=\(1-z\)%5E2-z%5E3+%3D+0](http://www.wolframalpha.com/input/?i=(1-z)%5E2-z%5E3+%3D+0)

from which we see that we have a small real and two larger complex conjugated singularities. We evaluate them numerically and see that their magnitude are  $\beta \approx 1.32471$  and  $\alpha \approx 0.5698402909980533$ . The dominant singularity is  $\alpha$ . We decomposed the denominator of  $S(z)$  into three roots of degree one. The function  $S(z)(z - \alpha)$  is therefore analytical at  $\alpha$ . This means that  $\alpha$  is a pole of first order.

For a meromorph generating function  $B(z)$  with poles  $\alpha_1, \dots, \alpha_m$  we know that there exist polynomials  $P_1(n), \dots, P_m(n)$  such that

$$[z^n]B(z) = \sum_{j=1}^m P_j(n)\alpha_j^n$$

and the degree of the polynomial  $P_j(n)$  is one smaller than the order of the pole  $\alpha_j$ . Since we have three poles of first order the polynomials are constants. We have

$$s_n = c_1(\alpha^{-n}) + c_2(\beta^{-n}).$$

for some constants  $c_1$  and  $c_2$ .

3. The exponential growth is  $\alpha^{-n} \approx 1.754877666246692655^n$ .
4. Notice that  $c_2(\beta^{-n}) = O(0.8^n)$ . If we can find the hidden factor  $c_1$  we get a good approximation with vanishing additive error for  $s_n$ . We can decompose

$$S(z) = \sum_{n=1}^{\infty} c_1 \alpha^{-n} z^n + c_2 \beta^{-n} z^n = \frac{c_1}{1 - z/\alpha} + \frac{c_2}{1 - z/\beta}$$

Let

$$B(z) = \frac{1}{1 - z/\alpha} \quad \text{and} \quad E(z) = \frac{1}{1 - z/\beta}$$

Notice that  $\lim_{z \rightarrow \alpha} B(z) = \infty$  while  $\lim_{z \rightarrow \alpha} E(z)$  is constant. Then

$$\lim_{z \rightarrow \alpha} \frac{S(z)}{B(z)} = \lim_{z \rightarrow \alpha} \frac{c_1 B(z) + c_2 E(z)}{B(z)} = c_1.$$

We use maxima to approximate

$$\lim_{z \rightarrow \alpha} \frac{S(z)}{B(z)} = \lim_{z \rightarrow \alpha} \frac{z(1-z)(1-z/\alpha)}{(1-z)^2 - z^3} \approx 0.23448675.$$

We use the following programm to verify the correctness

```

s = range(0,1000)
a = range(0,1000)
s[0] = 0
s[1] = 1
a[0] = 0
a[1] = 0
for n in range(2, 100):
    s[n] = a[n-2]+s[n-1]
    a[n] = a[n-1]+s[n-1]
    print n, s[n], 1.0*s[n]/s[n-1],
    0.23448675*1.754877666246692655**n/s[n]

```

The last line states

```
99 355268071453933228439241 1.75487766625 0.999999931817.
```

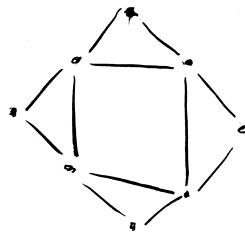
Indeed, after 100 iterations we only make a multiplicative error of 0.999999931817.

### Exercise T21

An algorithm  $I$  computes an optimal independent set for an undirected graph  $G = (V, E)$  of size  $n$  as follows: It picks a vertex  $v$  with maximal degree. If this degree is at most two, then the graph is a collection of cycles and paths and the solution is computed in linear time.

Otherwise, the optimal independent set either contains  $v$  (and then cannot contain any vertex in  $N(v)$ ) or it does not. Hence, the algorithm recursively computes the two independent sets  $I(G[V - N(v)])$  and  $I(G[V - \{v\}])$  and then chooses the bigger one, or the first if they have the same size.

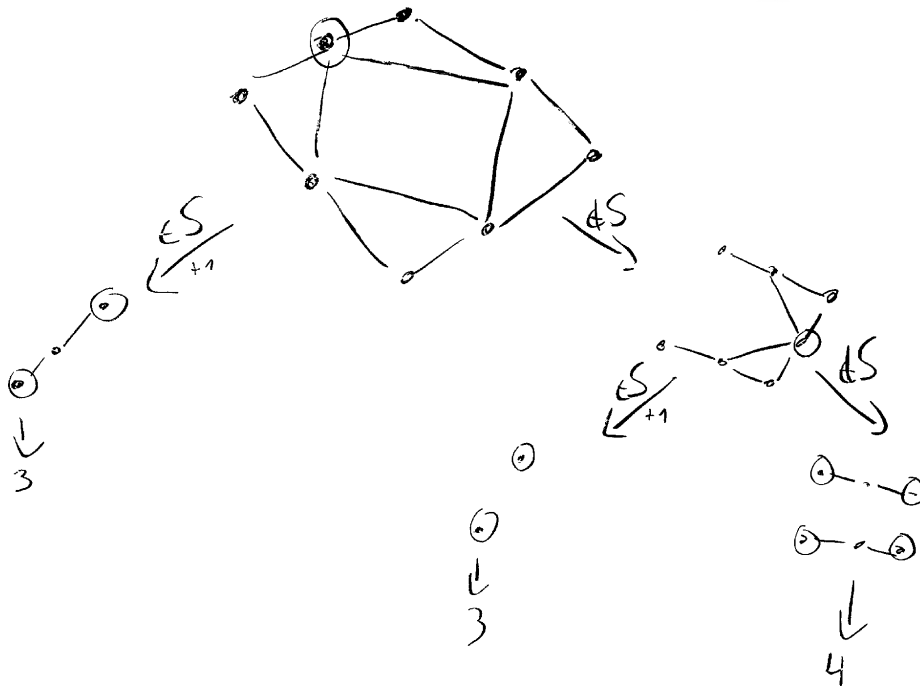
1. Simulate the algorithm on this graph:



2. Estimate its asymptotic running time up to a constant factor.

**Solution:**

1.



The best solution is therefore obtained by picking the 4 corner points.

2.

The smallest reduction happens if the graph has max degree 3, in this case we get the following recurrence relation:

$$t(n) = t(n - 4) + t(n - 1) + (n = 0)$$

We find a generating function:

$$\begin{aligned} S(z) &= z^4 S(z) + z S(z) + 1 \\ (1 - z^4 - z) S(z) &= 1 \\ S(z) &= \frac{1}{1 - z^4 - z} \end{aligned}$$

Finding the dominant singularity  $\alpha$  of  $S(z)$  by setting  $1 - z^4 - z = 0$ , yields the numerical result  $\alpha \approx 0.72449$ . And a runtime of the algorithm of  $O(\alpha^{-n}) = O(1.38028^n)$ .

### Exercise H14

In this exercise we will look at 2-3-trees. They are rooted, ordered trees. Each internal node has either two or three children. As usual, the size of a 2-3-tree will be the number of its internal nodes.

1. How can you define 2-3-trees recursively?
2. Enumerate all 2-3-trees of size two. How many are there? How many trees exist of sizes zero and one?
3. Find a generating function  $Q(z)$  for the number  $q_n$  of 2-3-trees with size  $n$ .
4. What is the dominant singularity of  $Q(z)$  and what is the exponential growth of  $q_n$ ? Use a computer algebra system. Do not give up when you see horrifying formulas.

**Solution:**

We get immediately the equation  $Q(z) = zQ(z)^2 + zQ(z)^3 + 1$ . If we solve it with the help of a computer algebra system, we get a very messy result, but we can spot the expression  $\sqrt{z^2 + 11z - 1}$ , which defines one of the singularities. It seems that it is probably the dominant one. If that is true, then  $\alpha = (5^{3/2} - 11)/2$  is the singularity for which we are looking. The exponential growth is then  $\alpha^{-n} \approx 11.09016994374933^n$ .

A closer look shows that this is indeed the dominant singularity.