

## Analysis of Algorithms

Recall from class that if a complex-valued function  $f(z)$  is analytic at 0 and  $R$  is the absolute value of the singularity nearest to the origin then the coefficient  $f_n = [z^n]f(z)$  satisfies:

$$f_n \asymp \left(\frac{1}{R}\right)^n.$$

### Exercise 11-1

Sort the following generating functions *within one minute* by their exponential growth!

1.  $A(z) = 1/\sqrt{1-z/2}$
2.  $B(z) = (1+z)/(1-z)$
3.  $C(z) = \frac{1}{1-e^{-z-1/3}}$

**Solution:** We only need to sort them by the absolute value of the dominant singularities.  $A_n \asymp 1/2^n$ ,  $B_n \asymp 1$ ,  $C_n \asymp 3^n$ .

### Exercise 11-2

An algorithm is given an array of length  $n \geq 0$  and, if  $n \geq 2$ , for each  $1 \leq k \leq n$  calls itself on some random subarray of length  $k$  with probability  $\frac{1}{2}$ . Compute the exponential growth of the running time of this algorithm.

**Solution:**

Let  $t_n$  be the expected number of calls of the algorithm for an array of length  $n$ . Then  $t_0 = t_1 = 1$  and  $t_n = 1 + \frac{1}{2} \sum_{k=1}^n t_k$  for  $n > 1$ . To get an equation that holds for all  $n$ , we let

$$t_n = 1 + \frac{1}{2} \sum_{k=1}^n t_k - \frac{1}{2}(n=1).$$

We make yet another modification so that in the equation above so that we may sum from  $k=0$ :

$$t_n = 1 + \frac{1}{2} \sum_{k=0}^n t_k - \frac{1}{2}(n=1) - \frac{1}{2}.$$

Define  $T(z) = \sum_{n=0}^{\infty} t_n z^n$ . Multiplying both sides by  $z^n$  and summing over  $n$ , we obtain:

$$T(z) = \frac{1}{1-z} + \frac{1}{2(1-z)}T(z) - \frac{1}{2}z + \frac{1}{2(1-z)}$$

which gives the following functional relation:  $T(z) = \frac{z^2-z+3}{1-2z}$ . Since the dominant singularity is located at  $z = \frac{1}{2}$ , we get an asymptotic running time of  $\asymp 2^n$ .

### Exercise 11-3

Determine  $[z^n]_{2-e^z}^{-1}$  up to an additive error of  $O(12^{-n})!$

**Solution:** We first determine the dominant singularity, which is  $\ln 2$ . Since, by L'Hôpital,

$$\lim_{z \rightarrow \ln 2} \frac{\ln 2 - z}{2 - e^z} = \frac{1}{2},$$

we have

$$\frac{\ln 2 - z}{\ln 2 - z} \cdot \frac{1}{2 - e^z} \sim \frac{1}{2} \cdot \frac{1}{\ln 2 - z} = \frac{1}{2 \ln 2} \cdot \frac{1}{1 - (1/\ln 2)z} = \frac{1}{2 \ln 2} \sum_{n=0}^{\infty} \left(\frac{1}{\ln 2}\right)^n z^n.$$

We now look for the next singularities (ordered by their distance from the origin). These are  $\ln 2 \pm 2\pi i$ . By L'Hôpital

$$\lim_{z \rightarrow \ln 2 \pm 2\pi i} \frac{z - \ln 2 \mp 2\pi i}{2 - e^z} = -\frac{1}{2}$$

and therefore

$$S(z) \sim -\frac{1}{2} \frac{1}{z - \ln 2 \mp 2\pi i} = \frac{1}{2} \frac{1}{\ln 2 \mp 2\pi i} \frac{1}{1 - \frac{z}{\ln 2 \mp 2\pi i}}$$

for  $z \rightarrow \ln 2 \pm 2\pi i$ . We can now use the Theorem of the lecture to determine the coefficients:

$$\begin{aligned} [z^n]S(z) &= \frac{1}{2} \left( \left(\frac{1}{\ln 2}\right)^{n+1} + \left(\frac{1}{\ln 2 + 2\pi i}\right)^{n+1} + \left(\frac{1}{\ln 2 - 2\pi i}\right)^{n+1} \right) + O(r)^{-n} \\ &= \frac{1}{2} \left( \left(\frac{1}{\ln 2}\right)^{n+1} + r^{n+1} (e^{i\phi(n+1)} + e^{-i\phi(n+1)}) \right) + O(r)^{-n} \\ &= \frac{1}{2} \left( \left(\frac{1}{\ln 2}\right)^{n+1} + r^{n+1} \cos(\phi(n+1)) \right) + O(r)^{-n} \end{aligned}$$

with  $r = 1/\sqrt{\ln^2 2 + 4\pi^2} \approx 12.58547409739904$ ,  $\phi = \arctan(\frac{2\pi}{\ln 2})$ .

### Homework Assignment 11-1 (10 points)

Determine the exponential growth of  $[z^n]G(z)$ , where

- $G(z) = z^2/(1 - z - z^2)$ ,
- $G(z) = \sqrt{1 + 2z} - \sqrt{2 + 2z - 4z^2}/\sqrt{3}$ ,
- $G(z) = \ln(1 + \sin(z))/\ln(1 + \cos(z))$ .

### Homework Assignment 11-2 (10 points)

$$A(z) = \frac{\sqrt{1 - z^7}}{2z^2 - 3z + 1} \quad B(z) = \frac{1 - z^2}{e^{z+3z^2}} \quad C(z) = z^5 + 3z^2(z^3 + z^2 + 8)$$

Sort the sequences  $a_n$ ,  $b_n$ , and  $c_n$  by their exponential growth.

### Homework Assignment 11-3 (10 Points)

Determine  $g_n = [z^n]G(z)$  up to an additive error of  $O(4^n)$ , where

$$G(z) = \sum_{n=0}^{\infty} g_n z^n = \frac{15z^2 + 8z + 1}{15z^2 - 8z + 1}.$$