

## Exercise Sheet 02

### Problem T3

In this warmup exercise we are going to analyse  $C'$ —the number of comparisons done in quick-sort during partitioning phases *on the left side* of the array. All comparisons are distributed among the following two lines in the program. We split the number of comparisons into  $C'$  for the first line and  $C''$  for the second line.

```
do { i++; } while(a[i] < k);  
do { j--; } while(k < a[j]);
```

We would expect  $C'$  should be around half the number of comparisons, i.e., around  $C/2$ .

### Problem T4

If a flow diagram consists of  $n$  nodes and  $m$  edges, how many fundamental cycles do we get?

### Problem T5

Prove or disprove: In every flow diagram you can find a spanning tree such that all fundamental cycles contain only edges that are labeled with plus.

### Problem T6

Let  $w \in \{a, b\}^n$  a word that has been chosen uniformly at random. How often is the body of the `while`-loop executed on average in the following algorithm? The function `is_palindrome` tests whether a word is a palindrome, i.e., the same when read backwards.

```
i = 2;  
while (i <= n) {  
    if (is_palindrome(w[1], ..., w[i])) return true;  
    i++;  
}  
return false;
```

### Problem H4 (15 credits)

We consider the following Algorithm. The array `a` contains a random permutation of the the numbers  $1, \dots, N$ .

```
void doSomething(int *a, int N) |  
{  
    int i; |  
  
    for (i=0; i<N-1; i++) /* 1 */  
        while (a[i] > a[i+1]) /* 2 */  
            a[i]--; /* 3 */  
}
```

How often is line 3 executed on average?

**Problem H5** (15 credits)

In this exercise, we consider Prim's Algorithm, which computes a minimum spanning tree. The input to this algorithm is a graph  $G = (V, E)$ , a weight function on the edges  $w: E \rightarrow \mathbf{R}$  and a starting node  $r$ .

```
1  for each  $u \in V$  do
2       $key[u] \leftarrow \infty$ 
3       $\pi[u] \leftarrow \text{NIL}$ 
4   $key[r] \leftarrow 0$ 
5   $M \leftarrow V$ 
6  while ( $M \neq \emptyset$ ) do
7       $u \leftarrow \text{min-from}(M)$ 
8      for each  $v \in \text{neighbors}(u)$  do
9          if ( $v \in M$ )  $\wedge$  ( $w(u, v) < key[v]$ ) then
10              $\pi[v] \leftarrow u$ 
11              $key[v] \leftarrow w(u, v)$ 
```

Construct the control flow graph, a spanning tree in the control flow graph, the fundamental cycles, a corresponding linear system of equations and a solution to this system.