

Exercise Sheet 02

Problem T3

In this warmup exercise we are going to analyse C' —the number of comparisons done in quick-sort during partitioning phases *on the left side* of the array. All comparisons are distributed among the following two lines in the program. We split the number of comparisons into C' for the first line and C'' for the second line.

```
do { i++; } while(a[i] < k);  
do { j--; } while(k < a[j]);
```

We would expect C' should be around half the number of comparisons, i.e., around $C/2$.

Problem T4

If a flow diagram consists of n nodes and m edges, how many fundamental cycles do we get?

Problem T5

Prove or disprove: In every flow diagram you can find a spanning tree such that all fundamental cycles contain only edges that are labeled with plus.

Problem T6

Let $w \in \{a, b\}^n$ a word that has been chosen uniformly at random. How often is the body of the `while`-loop executed on average in the following algorithm? The function `is_palindrome` tests whether a word is a palindrome, i.e., the same when read backwards.

```
i = 2;  
while (i <= n) {  
    if (is_palindrome(w[1], ..., w[i])) return true;  
    i++;  
}  
return false;
```

Problem H4 (15 credits)

We consider the following Algorithm. The array `a` contains a random permutation of the the numbers $1, \dots, N$.

```
void doSomething(int *a, int N) |  
{  
    int i; |  
  
    for (i=0; i<N-1; i++) /* 1 */  
        while (a[i] > a[i+1]) /* 2 */  
            a[i]--; /* 3 */  
}
```

How often is line 3 executed on average?

Problem H5 (15 credits)

In this exercise, we consider Prim's Algorithm, which computes a minimum spanning tree. The input to this algorithm is a graph $G = (V, E)$, a weight function on the edges $w: E \rightarrow \mathbf{R}$ and a starting node r .

```
1  for each  $u \in V$  do
2       $key[u] \leftarrow \infty$ 
3       $\pi[u] \leftarrow \text{NIL}$ 
4   $key[r] \leftarrow 0$ 
5   $M \leftarrow V$ 
6  while ( $M \neq \emptyset$ ) do
7       $u \leftarrow \text{min-from}(M)$ 
8      for each  $v \in \text{neighbors}(u)$  do
9          if ( $v \in M$ )  $\wedge$  ( $w(u, v) < key[v]$ ) then
10              $\pi[v] \leftarrow u$ 
11              $key[v] \leftarrow w(u, v)$ 
```

Construct the control flow graph, a spanning tree in the control flow graph, the fundamental cycles, a corresponding linear system of equations and a solution to this system.